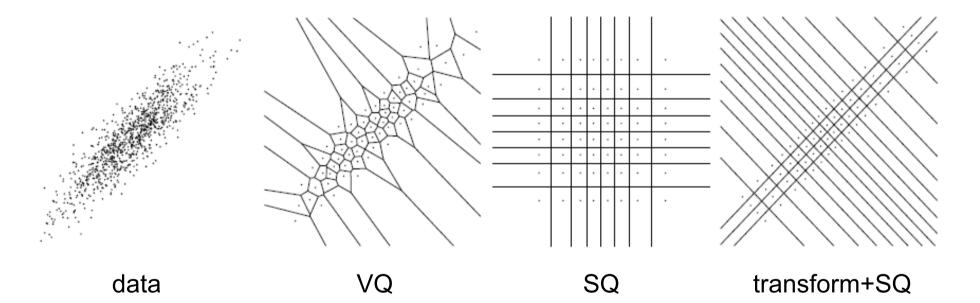
Transform Coding

- Coding sources with memory
- Orthonormal transforms
- Rate distortion function for transform coding
- Transform coding gain
- Maximizing the transform coding gain
- Karhunen-Loeve Transform (KLT)
- Energy concentration by KLT



Coding Sources with Memory

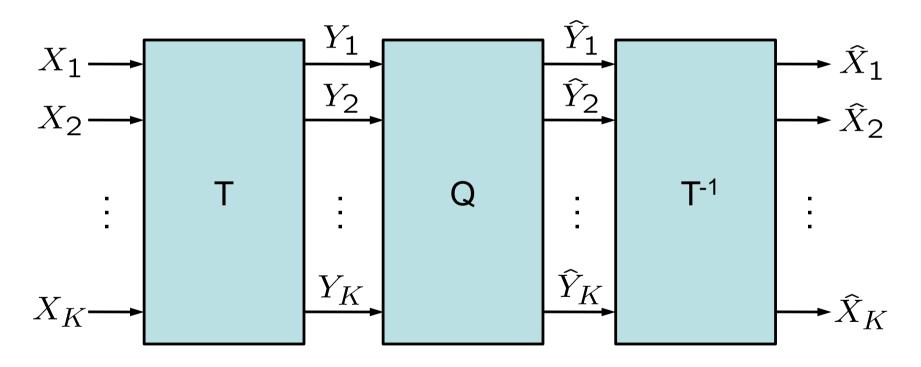
Example: Two-dimensional source



Exploit memory of the source by linear transform

Using Transforms for Coding Vector Sources

K-dimensional vector source

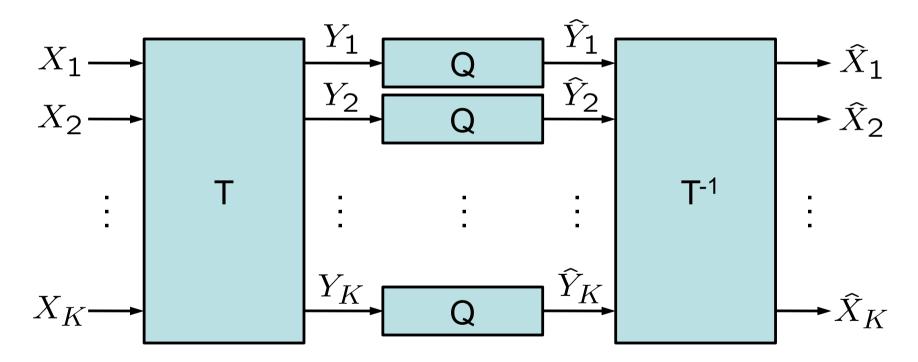


• How to choose the transform *T*?



Independent Transform Coefficients

Let the transform generate independent coefficients



Further, use independent quantizers (no sphere packing)



Orthonormal Transforms

- Consider a transform T that produces independent transform coefficients
- The synthesis transform exists and is given by T⁻¹
- The component quantizers reconstruct to their respective centroids
- **Theorem:** To minimize the mean square error distortion, it is sufficient to consider transforms that satisfy $TT^{T} = I$. [Goyal, 2001]



Energy Conservation

For any orthonormal transform y= Tx we obtain

$$||y||_2^2 = y^T y = x^T T^T T x = ||x||_2^2$$

- Vector lengths ("energies") are conserved.
- Interpretation: every orthonormal transform is simply a rotation of the coordinate system.



Energy Distribution for Orthonormal Transform

- Energy is conserved, but typically will be unevenly distributed among coefficients.
- Autocorrelation matrix

$$C_{yy} = E\{yy^T\} = E\{Txx^TT^T\} = TC_{xx}T^T$$

Mean squared values ("average energies") of the coefficients y_i are on the diagonal of C_{yy}



Distortion Rate Function for Transform Coding

Component distortion at high rates

$$D_n(R_n) = \epsilon^2 \sigma_{y_n}^2 2^{-2R_n}$$
 for $D_n \le \epsilon^2 \sigma_{y_n}^2$

Optimum rate allocation

$$\min \frac{1}{N} \sum_{n=1}^{N} D_n(R_n)$$
 s.t. $R = \frac{1}{N} \sum_{n=1}^{N} R_n$

Optimum rate allocation at high rates

$$D_n \le \epsilon^2 \sigma_{y_n}^2 \quad : \quad \frac{\partial (D + \lambda R)}{\partial R_i} = 0 \quad \Rightarrow \quad \lambda = 2 \ln(2) D_i$$
$$D_i = D_j \quad \forall \quad i, j$$



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Transform Coding no. 8

Distortion Rate Function for Transform Coding

Component rate at high rates

$$R_n = \frac{1}{2} \log_2 \left(\frac{\epsilon^2 \sigma_{y_n}^2}{D_n} \right) \quad \text{for} \quad D_n \le \epsilon^2 \sigma_{y_n}^2$$

Total rate at high rates

$$R = \frac{1}{N} \sum_{n=1}^{N} \frac{1}{2} \log_2 \left(\frac{\epsilon^2 \sigma_{y_n}^2}{D} \right) = \frac{1}{2} \log_2 \left(\frac{\epsilon^2 \left[\prod_{n=1}^{N} \sigma_{y_n}^2 \right]^{\frac{1}{N}}}{D} \right)$$

Distortion rate function for transform coding

$$D(R) = \epsilon^2 \left[\prod_{n=1}^N \sigma_{y_n}^2 \right]^{\frac{1}{N}} 2^{-2R} \quad \text{for} \quad D \le \epsilon^2 \sigma_{y_n}^2$$



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Transform Coding no. 9

Transform Coding Gain

- Distortion gain of transform coding with block size N over scalar quantization of source samples.
- Transform coding gain

$$G_{TC} = \frac{D_{SQ}(R)}{D_{TC}(R)}$$

$$G_{TC} = \frac{\sigma_x^2}{\left[\prod_{n=1}^N \sigma_{y_n}^2\right]^{\frac{1}{N}}} = \frac{\frac{1}{N} \sum_{n=1}^N \sigma_{y_n}^2}{\left[\prod_{n=1}^N \sigma_{y_n}^2\right]^{\frac{1}{N}}}$$



Remark: Hadamard Inequality

- Consider a symmetric positive-definite matrix A
- Choose B, b, and x such that Bx = -b

$$\det \begin{bmatrix} B & b \\ b^{H} & c \end{bmatrix} = \det \begin{bmatrix} I & 0 \\ x^{H} & 1 \end{bmatrix} \begin{bmatrix} B & b \\ b^{H} & c \end{bmatrix} \begin{bmatrix} I & x \\ 0 & 1 \end{bmatrix}$$
$$= \det \begin{bmatrix} B & 0 \\ 0 & b^{H}x + c \end{bmatrix}$$
$$= \det(B)(b^{H}x + c)$$
$$\leq \det(B)c$$

• Note that $b^H x = -x^H B x$ is negative

$$\det(A) \le \prod_i A_{ii}$$

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Transform Coding no. 11

Maximizing Transform Coding Gain

Lower bound for product of coefficient variances

$$\det(C_{yy}) \leq \prod_{n=1}^N \sigma_{y_n}^2$$

Maximize transform coding gain

$$\max G_{TC} \quad \Leftrightarrow \quad \min \prod_{n=1}^{N} \sigma_{y_n}^2$$

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Transform with maximum transform coding gain satisfies

$$\det(C_{yy}) = \prod_{n=1}^{N} \sigma_{y_n}^2$$

The transform that decorrelates the signal maximizes the transform coding gain

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Eigenmatrix of the Autocorrelation Matrix

Definition: Eigenmatrix Φ of the autocorrelation matrix C_{xx}

- $-\Phi$ is unitary
- The columns of Φ form a set of eigenvectors of C_{xx} , i.e.,

$$C_{xx} \Phi = \Phi \Lambda \qquad \qquad \Lambda = \begin{bmatrix} \lambda_1 & & 0 \\ & \lambda_2 & \\ & & \ddots & \\ 0 & & & \lambda_K \end{bmatrix}$$

$$\Lambda \text{ is a diagonal matrix} \qquad \qquad 0 \qquad \qquad 0 \qquad \qquad 0$$

- − C_{xx} is symmetric nonnegative definite, hence $\lambda_i \ge 0$ for all i
- C_{xx} is normal matrix, i.e., $C_{xx}^{H} C_{xx} = C_{xx} C_{xx}^{H}$, hence unitary eigenmatrix exists



Karhunen-Loeve Transform

Unitary transform with matrix

$$T = \Phi^H$$

where the columns of Φ are ordered according to decreasing eigenvalues.

Transform coefficients are pairwise uncorrelated

$$C_{yy} = TC_{xx}T^H = \Phi^H C_{xx}\Phi = \Phi^H \Phi \Lambda = \Lambda$$

The KLT maximizes the transform coding gain.

Energy Concentration Property of KLT

- No other unitary transform packs as much energy into the first J coefficients, where J is arbitrary
- Mean squared approximation error by choosing only first J coefficients is minimized.



Optimum Energy Concentration by KLT

 To show optimum energy concentration property, consider the truncated coefficient vector

$$b = I_J y$$

- Where I_J contains ones on the first J diagonal positions, else zeros.
- Energy in first J coefficients for arbitrary transform A

$$E = tr(C_{bb}) = tr(I_J C_{yy} I_J) = tr(I_J T C_{xx} T^H I_J) = \sum_{k=1}^J t_k^T C_{xx} t_k^*$$

T

where t_k^{T} is the k-th row of T.

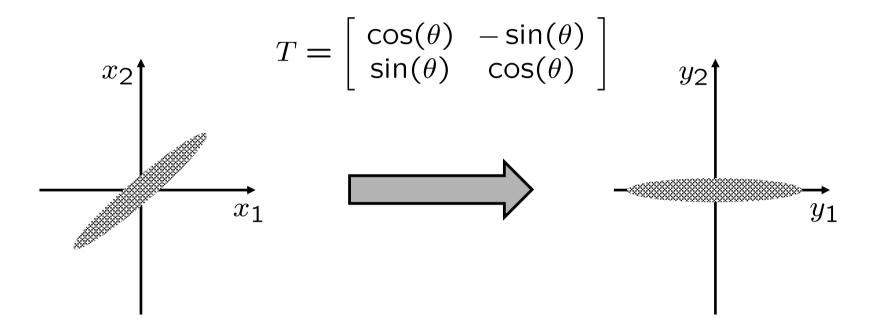
- Minimize Lagrangian cost function to enforce unit-length basis vectors $L = -E + \sum_{k=1}^{J} \lambda_k (t_k^T t_k^* - 1) = -\sum_{k=1}^{J} t_k^T C_{xx} t_k^* + \sum_{k=1}^{J} \lambda_k (t_k^T t_k^* - 1)$
- Differentiating L with respect to t_j yields necessary condition

$$C_{xx}t_j^* = \lambda_j t_j^* \quad \forall \quad j \le J$$



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Illustration of Energy Concentration



Before KLT: Strongly correlated samples, equal energies After KLT: Uncorrelated samples, most of the energy in first coefficient



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