Rate Distortion Theory

- Distortion measure
- Rate distortion function
- Shannon's noisy source coding theorem
- Convexity of rate distortion function
- Shannon lower bound
- Rate distortion function for Gaussian source



Encoding and Decoding

Encoding and decoding as data processing



- Lossy coding: $X \approx \hat{X}$
- Data processing inequality: $I(X; \hat{X}) \leq I(X; Y)$

Encoding and Decoding



 $I(X; \hat{X}) \leq I(X; Y) \leq H(Y)$



Rate Distortion Theory

 Rate distortion theory calculates the minimum transmission bitrate R for a required reconstruction quality



 Results of rate distortion theory are obtained without consideration of a specific coding method

Distortion Measure

- Symbol (signal, image . . .) x sent, \hat{x} received
- Single-letter distortion measure:

$$d(x, \hat{x}) \geq 0$$

 $d(x, \hat{x}) = 0$ for $x = \hat{x}$

Average distortion:

$$E\{d(X,\hat{X})\} = \sum_{x} \sum_{\hat{x}} f_{X,\hat{X}}(x,\hat{x})d(x,\hat{x})$$

• Distortion criterion: $E\{d(X, \hat{X})\} \leq D$

Maximum permissible average distortion

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Rate Distortion Function

 Lower the bit-rate R by allowing some acceptable distortion D of the signal





Rate Distortion Function

Definition:

$$R(D) = \inf_{f_{\widehat{X}|X}} I(X; \widehat{X}) \quad \text{s.t.} \quad E\{d(X, \widehat{X})\} \le D$$

Shannon's Noisy Source Coding Theorem:

For a given maximum average distortion D, the rate distortion function R(D) is the (achievable) lower bound for the transmission bit-rate.

- R(D) is continuous, monotonically decreasing for R>0 and convex
- Equivalently use distortion-rate function D(R)



Convexity of the Rate Distortion Function

- Consider two rate distortion pairs (R_a, D_a) and (R_b, D_b) which lie on the rate distortion curve.
- Let the joint distributions $a_{X,\hat{X}}$ and $b_{X,\hat{X}}$ achieve these pairs.
- Consider the distribution: $f_{X,\widehat{X}} = \lambda a_{X,\widehat{X}} + [1 \lambda]b_{X,\widehat{X}}$





Remark: Convexity of Mutual Information

- Mutual information is a convex function of the conditional distribution
- Consider the distribution: $f_{\widehat{X}|X} = \lambda a_{\widehat{X}|X} + [1 \lambda]b_{\widehat{X}|X}$





Remark: Convexity of Mutual Information

• For the conditional distribution $f_{\widehat{X}|X}$ we have

$$\begin{split} I_f(X;\hat{X}) &= \sum f_X f_{\hat{X}} \frac{f_{\hat{X}|X}}{f_{\hat{X}}} \log_2 \frac{f_{\hat{X}|X}}{f_{\hat{X}}} \\ &= \sum f_X f_{\hat{X}} \left[\lambda \frac{a}{f_{\hat{X}}} + [1-\lambda] \frac{b}{f_{\hat{X}}} \right] \log_2 \left[\lambda \frac{a}{f_{\hat{X}}} + [1-\lambda] \frac{b}{f_{\hat{X}}} \right] \\ &\leq \sum f_X f_{\hat{X}} \left[\lambda \frac{a}{f_{\hat{X}}} \log_2 \frac{a}{f_{\hat{X}}} + [1-\lambda] \frac{b}{f_{\hat{X}}} \log_2 \frac{b}{f_{\hat{X}}} \right] \\ &= \lambda I_a(X;\hat{X}) + [1-\lambda] I_b(X;\hat{X}) \end{split}$$

• Note that $g(t) = t \log_2 t$ is convex, i.e.,

 $g(\lambda t_1 + [1 - \lambda]t_2) \le \lambda g(t_1) + [1 - \lambda]g(t_2)$



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Convexity of the Rate Distortion Function

• For the joint distribution $f_{X,\hat{X}}$ we have

$$R_f(D_f) = \inf_{f_{\widehat{X}|X}} I_f(X; \widehat{X}) \quad \text{s.t.} \quad E_f\{d(X, \widehat{X})\} \le D$$

- Distribution of the source is given: f_X
- Mutual information is a convex function of $f_{\widehat{X}|X}$

$$R_f(D_f) = \inf I_f(X; \hat{X})$$

$$\leq \inf \left\{ \lambda I_a(X; \hat{X}) + [1 - \lambda] I_b(X; \hat{X}) \right\}$$

$$= \lambda \inf I_a(X; \hat{X}) + [1 - \lambda] \inf I_b(X; \hat{X})$$

$$= \lambda R_a(D_a) + [1 - \lambda] R_b(D_b)$$



Characterization of the Rate Distortion Function

• Solve unconstrained problem $\inf_{f_{\widehat{X}|X}} J(X; \widehat{X})$ with $\lambda > 0$ and

$$J(X;\hat{X}) = E\left\{ \ln \frac{f_{\hat{X}|X}(\hat{X}|X)}{f_{\hat{X}}(\hat{X})} + \lambda[d(X,\hat{X}) - D] \right\}$$
$$= E\left\{ -\ln \left[\frac{f_{\hat{X}}(\hat{X})}{f_{\hat{X}|X}(\hat{X}|X)} e^{-\lambda[d(X,\hat{X}) - D]} \right] \right\}$$
$$\geq -\ln E\left\{ \frac{f_{\hat{X}}(\hat{X})}{f_{\hat{X}|X}(\hat{X}|X)} e^{-\lambda[d(X,\hat{X}) - D]} \right\}$$
$$= -\ln E\{g(\hat{X}|X)\}$$



Characterization of the Rate Distortion Function

- Jensen: Lower bound is tight iff $g(\hat{x}|x)$ is constant in \hat{x}
- We obtain the conditional PDF

$$f_{\widehat{X}|X}(\widehat{x}|x) = \frac{f_{\widehat{X}}(\widehat{x})e^{-\lambda d(x,\widehat{x})}}{\mu(x)}$$

Normalizing the conditional PDF

$$f_{\widehat{X}|X}(\widehat{x}|x) = \frac{f_{\widehat{X}}(\widehat{x})e^{-\lambda d(x,\widehat{x})}}{\int f_{\widehat{X}}(\xi)e^{-\lambda d(x,\xi)}d\xi}$$

 Basis for Blahut algorithm, which is an iterative method to find a numerical solution for the rate distortion function

Shannon Lower Bound

- It can be shown that $h(X \hat{X} | \hat{X}) = h(X | \hat{X})$
- Thus $R(D) = \inf_{d \le D} \{h(X) - h(X \mid \hat{X})\}$ $= h(X) - \sup_{d \le D} \{h(X \mid \hat{X})\}$ $= h(X) - \sup_{d \le D} \{h(X - \hat{X} \mid \hat{X})\}$
- Ideally, the source coder would introduce errors $x \hat{x}$ that are statistically independent from the reconstructed signal \hat{x} (not always possible!).
- Shannon lower bound:

$$R(D) \ge h(X) - \sup_{d \le D} h(X - \hat{X})$$



Shannon Lower Bound

 Mean squared error distortion measure: Gaussian PDF possesses largest entropy for given variance

$$R(D) \ge h(X) - \sup_{d \le D} h(X - \hat{X})$$
$$= h(X) - \frac{1}{2}\log_2 2\pi eD$$

Distortion reduction by 6 dB requires 1 bit/sample



R(*D*) Function for a Memoryless Gaussian Source and MSE Distortion

- Gaussian source, variance σ^2
- Mean squared error

 $E\{(X-\widehat{X})^2\} \le D$

- Rule of thumb: $6 dB \cong 1$ bit
- *R(D)* for non-Gaussian sources with the same variance σ² is always <u>below</u> this Gaussian *R(D)* curve.





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R(D) for Gaussian Source with Memory

- Bandlimited, jointly Gaussian source with power spectral density $\Phi_{xx}(\omega)$
- Mean squared error distortion $E\{(X \hat{X})^2\} \le D$
- Parametric formulation of the R(D) function

$$D(\theta) = \frac{1}{2\pi} \int_{\omega} \min\left\{\theta, \Phi_{xx}(\omega)\right\} d\omega$$
$$R(\theta) = \frac{1}{2\pi} \int_{\omega} \max\left\{0, \frac{1}{2}\log\frac{\Phi_{xx}(\omega)}{\theta}\right\} d\omega$$

R(*D*) for non-Gaussian sources with the same power spectral density is always lower.



