Continuous Random Variables

- Differential entropy
- Maximum entropy distributions
- Normal distribution
- Multivariate normal distribution
- Gaussian process
- Differential entropy rate



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Differential Entropy

 Let X be a continuous random variable with PDF f_X(x). The differential entropy of X is defined as

$$h(X) = E\left\{-\log_2 f_X\right\} = -\int_{\mathcal{R}} f_X(x)\log_2 f_X(x)dx$$

Example: Differential Entropy for uniform distribution

$$f_X(x) = \frac{1}{a} \mathbf{1}_{[0,a]}(x)$$
$$h(X) = \log_2(a)$$

Note, differential entropy can be negative!

Properties of Differential Entropy

Translation does not change the differential entropy

$$h(X+c) = h(X)$$

Scaling by a increases the differential entropy by log₂|a|

$$h(aX) = h(X) + \log_2|a|$$



Joint and Conditional Differential Entropy

Joint differential entropy of X and Y

$$h(X,Y) = E\left\{-\log_2 f_{X,Y}\right\}$$

- Conditional differential entropy of X given Y $h(X|Y) = E\left\{-\log_2 f_{X|Y}\right\}$ h(X|Y) = h(X,Y) h(Y)
- Conditioning reduces differential entropy $h(X|Y) \le h(X)$

with equality iff X and Y are independent.



Mutual Information

Mutual information between X and Y

$$I(X;Y) = E\left\{\log_2 \frac{f_{X,Y}}{f_X f_Y}\right\}$$
$$I(X;Y) = h(X) - h(X|Y)$$
$$I(X;Y) = h(Y) - h(Y|X)$$

Mutual information is non-negative

 $I(X;Y) \ge 0$

with equality iff X and Y are independent.



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Remark: Convex Functions

A function g(x) is said to be convex over an interval [a,b] if for every x₁, x₂ ∈ [a,b] and 0 ≤ λ ≤ 1



 A function is said to be strictly convex if equality holds only for λ = 0 or λ = 1.



Remark: Jensen's Inequality

If g is a convex function and X is a random variable, then

 $E\left\{g(X)\right\} \ge g\left(E\{X\}\right)$

- Moreover, if g is strictly convex, then equality implies that E{X}=X with probability 1, i.e., X is a constant.
- Example: Two mass point distribution $p_1g(x_1) + p_2g(x_2) \ge g(p_1x_1 + p_2x_2)$
- If g is strictly convex, equality holds only for the point x₁ = x₂ := x, that is, X is a constant.

Maximum Entropy Distribution for given σ^2

For given variance, find PDF that maximizes entropy

$$\inf_{f_X} \{-h(X)\} \quad \text{s.t.} \quad E\{X^2\} = \sigma^2$$

• Solve unconstrained problem $\inf_{f_X} J(X)$ with $\lambda > 0$ and

$$J(X) = E\left\{ \ln f_X(X) + \lambda(X^2 - \sigma^2) \right\}$$
$$= E\left\{ -\ln \frac{e^{-\lambda(X^2 - \sigma^2)}}{f_X(X)} \right\}$$
$$\ge -\ln E\left\{ \frac{e^{-\lambda(X^2 - \sigma^2)}}{f_X(X)} \right\}$$
$$= -\ln E\{Y\}$$



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Maximum Entropy Distribution for given σ^2

Jensen: Lower bound is tight iff Y is a constant.

$$\frac{e^{-\lambda(x^2-\sigma^2)}}{f_X(x)} = \text{const.} \quad \forall x \in \mathcal{R}$$

We obtain the PDF

$$f_X(x) = c e^{-\lambda x^2}$$

• Normal distribution with variance σ^2 that integrates to 1

$$f_X(x) = \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{x^2}{2\sigma^2}}$$



Differential Entropy of Normal Distribution

For given variance, maximum differential entropy is

$$h(X) = E \{ -\log_2 f_X \}$$

= $\frac{1}{2} \log_2(2\pi\sigma^2) + \frac{E\{X^2\}}{2\sigma^2} \log_2 e$
= $\frac{1}{2} \log_2(2\pi\sigma^2) + \frac{1}{2} \log_2 e$
= $\frac{1}{2} \log_2(2\pi e\sigma^2)$



Multivariate Normal Distribution

Let X=(X₁, X₂, ..., X_n)^T have a multivariate normal distribution with zero mean and covariance matrix C

$$f_{\mathbf{X}}(\mathbf{x}) = \frac{1}{(2\pi)^{\frac{n}{2}} |C|^{\frac{1}{2}}} e^{-\frac{1}{2}\mathbf{x}^{T}C^{-1}\mathbf{x}}$$

 Change of coordinate system: Y = T^T X, where T^T C T=Λ being diagonal with Eigenvalues λ_i; note that |T| = 1

$$\mathbf{x}^T C^{-1} \mathbf{x} = \mathbf{y}^T T^T C^{-1} T \mathbf{y} = \mathbf{y}^T \Lambda^{-1} \mathbf{y}$$
$$f_{\mathbf{Y}}(\mathbf{y}) = \prod_{i=1}^n \frac{1}{\sqrt{2\pi}\sqrt{\lambda_i}} e^{-\frac{1}{2}\frac{\mathbf{y}_i^2}{\lambda_i}}$$



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Multivariate Normal Distribution

 Differential entropy of a multivariate normal random variable

$$h(\mathbf{Y}) = h(T^T \mathbf{X})$$

= $h(\mathbf{X}) + \log_2 |T^T|$
= $h(\mathbf{X})$

$$h(\mathbf{X}) = \frac{n}{2} \log_2(2\pi e) + \sum_{i=1}^n \frac{1}{2} \log_2 \lambda_i$$

= $\frac{n}{2} \log_2(2\pi e) + \frac{1}{2} \log_2 |C|$
= $\frac{1}{2} \log_2[(2\pi e)^n |C|]$



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Bivariate Normal Random Variable

Covariance matrix C

$$C = \begin{bmatrix} \sigma_1^2 & \sigma_1 \sigma_2 \rho_{12} \\ \sigma_1 \sigma_2 \rho_{12} & \sigma_2^2 \end{bmatrix} \quad |\rho_{12}| \le 1$$

Mutual information between components

$$I(X_{1}; X_{2}) = h(X_{1}) + h(X_{2}) - h(X_{1}, X_{2})$$

= $\frac{1}{2} \log_{2} \frac{\left| \begin{bmatrix} \sigma_{1}^{2} & 0 \\ 0 & \sigma_{2}^{2} \end{bmatrix} \right|}{|C|}$
= $-\frac{1}{2} \log_{2}(1 - \rho_{12}^{2})$

•
$$I(X_1;X_2) = 0$$
 for $\rho_{12} = 0$



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Continuous-Valued Stochastic Process

Continuous-valued discrete-time stochastic process {X_i}, X_i 2 Rⁿ, is an indexed sequence of continuous random variables X₁, X₂, ..., X_n with joint PDF

$$f_{X_1, X_2, \dots, X_n} = \Pr \{ (X_1, X_2, \dots, X_n) = (x_1, x_2, \dots, x_n) \}$$

$$\forall \quad (x_1, x_2, \dots, x_n) \in \mathcal{R}^n$$

 A stochastic process is said to be stationary if the joint PDF is invariant with respect to shifts

$$f_{X_1,X_2,\dots,X_n} = f_{X_{1+l},X_{2+l},\dots,X_{n+l}} \quad \forall \quad l \in \mathcal{Z}$$

Differential Entropy Rate

- How does the differential entropy of Xⁿ grow with n?
- Differential entropy rate of a continuous-valued stochastic process {X_i} is defined by

$$h(\{X_i\}) = \lim_{n \to \infty} \frac{1}{n} h(X_1, X_2, \dots, X_n)$$

Using the chain rule

$$h(\{X_i\}) = \lim_{n \to \infty} \frac{1}{n} \sum_{i=1}^n h(X_i | X_{i-1}, \dots, X_1)$$



Gaussian Stochastic Process

 Consider a zero-mean cyclostationary Gaussian process
{X_i} whose statistical properties repeat with period n

$$h(\{X_i\}) = \frac{1}{n}h(X_1, X_2, \dots, X_n)$$

= $\frac{1}{2}\log_2(2\pi e) + \frac{1}{n2}\log_2|C^{(n)}|$

• Let the covariance matrix be **circulant** with $C_{kl} = E\{X_k X_l\}$

$$C^{(n)} = \begin{bmatrix} c_0 & c_1 & c_2 & \cdots & c_{n-1} \\ c_{n-1} & c_0 & c_1 & \cdots & c_{n-2} \\ c_{n-2} & c_{n-1} & c_0 & \cdots & c_{n-3} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ c_1 & c_2 & c_3 & \cdots & c_0 \end{bmatrix} \quad C_{kl} = c_{(l-k)} \mod n$$



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Remark: Eigenvalues of Circulant Matrix

• Let C have Eigenvalues λ_v and Eigenvectors $t^{(v)}$

$$Ct^{(\nu)} = \lambda_{\nu}t^{(\nu)}$$
$$\sum_{l=1}^{n} C_{kl}t_{l} = \lambda t_{k} \quad \forall \quad k = 1, \dots, n$$
$$\sum_{l=0}^{n-k} c_{l}t_{l+k} + \sum_{l=n-k+1}^{n-1} c_{l}t_{l+k-n} = \lambda t_{k} \quad \forall \quad k = 1, \dots, n$$

Solve difference equation by using complex roots of unity

$$\lambda_{\nu} = \sum_{l=0}^{n-1} c_l e^{-j\frac{2\pi}{n}l\nu} := \mathsf{DFT}\{c_l\}$$



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Gaussian Stochastic Process

 Consider a zero-mean cyclostationary Gaussian process with period n

$$h(\{X_i\}) = \frac{1}{2}\log_2(2\pi e) + \frac{1}{n}\frac{1}{2}\log_2|C^{(n)}|$$

= $\frac{1}{2}\log_2(2\pi e) + \frac{1}{n}\sum_{\nu=0}^{n-1}\frac{1}{2}\log_2\lambda_{\nu}$
= $\frac{1}{2}\log_2(2\pi e) + \frac{1}{n}\sum_{\nu=0}^{n-1}\frac{1}{2}\log_2\mathsf{DFT}\{c_l\}_{\nu}$



Remark: Szegö's Theorem

 Let the period n grow and consider a stationary Gaussian stochastic process with autocorrelation sequence and PSD

$$\phi_{XX}[l] = E\{X_i X_{i+l}\} \qquad \Phi_{XX}(\omega) = \sum_{l \in \mathcal{Z}} \phi_{XX}[l] e^{-j\omega l}$$

• Then, the covariance matrix is Toeplitz with Eigenvalues λ_{v}

$$C^{(n)} = \begin{bmatrix} \phi_0 & \phi_1 & \phi_2 & \cdots & \phi_{n-1} \\ \phi_{-1} & \phi_0 & \phi_1 & \cdots & \phi_{n-2} \\ \phi_{-2} & \phi_{-1} & \phi_0 & \cdots & \phi_{n-3} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \phi_{1-n} & \phi_{2-n} & \phi_{3-n} & \cdots & \phi_0 \end{bmatrix}$$

The distribution of discrete Eigenvalues converges

$$\lim_{n \to \infty} \frac{1}{n} \sum_{\nu=0}^{n-1} g(\lambda_{\nu}) = \frac{1}{2\pi} \int_{-\pi}^{\pi} g\left(\Phi(\omega)\right) d\omega$$



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Gaussian Stochastic Process

- Given a zero-mean stationary Gaussian process {X_i} with $\phi_{XX}[l] = E\{X_i X_{i+l}\}$ $\Phi_{XX}(\omega) = \sum \phi_{XX}[l] e^{-j\omega l}$
- The differential entropy rate is

$$h(\{X_i\}) = \frac{1}{2}\log_2(2\pi e) + \lim_{n \to \infty} \frac{1}{n} \sum_{\nu=0}^{n-1} \frac{1}{2}\log_2 \lambda_{\nu}$$

$$= \frac{1}{2}\log_2(2\pi e) + \lim_{n \to \infty} \frac{1}{n} \sum_{\nu=0}^{n-1} \frac{1}{2}\log_2 \mathsf{DFT}\{c_l\}_{\nu}$$

$$= \frac{1}{2}\log_2(2\pi e) + \frac{1}{2\pi} \int_{-\pi}^{\pi} \frac{1}{2}\log_2 \Phi_{XX}(\omega) d\omega$$



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Continuous RVs no. 20

 $l \in \mathcal{Z}$