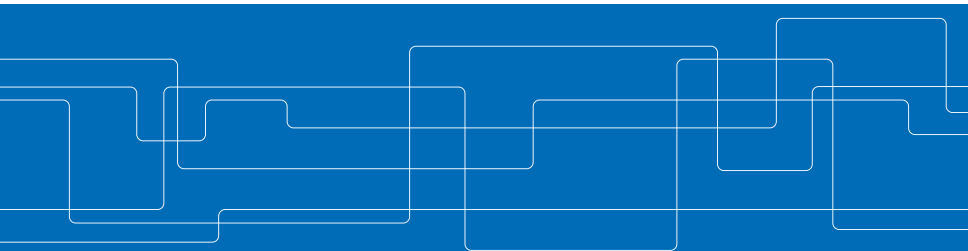




Secure and Privacy Preserving Vehicular Communication Systems: Identity and Credential Management Infrastructure

Mohammad Khodaei
Networked Systems Security Group (NSS)

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Vehicular Communication (VC) Systems

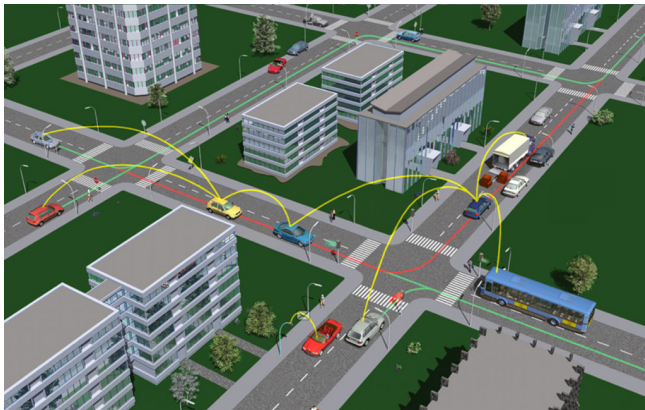
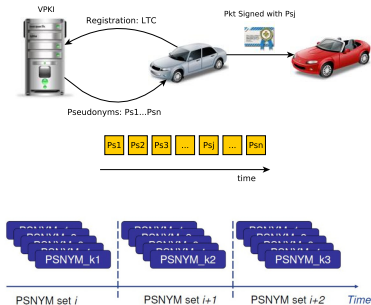


Figure: Photo Courtesy of the Car2Car Communication Consortium (C2C-CC)

Security and Privacy for VC Systems¹

Basic Requirements

- ▶ Message authentication & integrity
- ▶ Message non-repudiation
- ▶ Access control
- ▶ Entity authentication
- ▶ Accountability
- ▶ Privacy protection



Vehicular Public-Key Infrastructure (VPKI)

- ▶ Pseudonymous authentication
- ▶ Trusted Third Party (TTP):
 - ▶ Certification Authority (CA)
 - ▶ Issues credentials & binds users to their pseudonyms

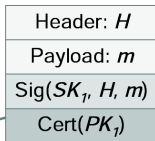
¹ P. Papadimitratos, et al. "Securing Vehicular Communications - Assumptions, Requirements, and Principles," in ESCAR, Berlin, Germany, pp. 5-14, Nov. 2006.

P. Papadimitratos, et al. "Secure Vehicular Communication Systems: Design and Architecture," in IEEE Communications Magazine, vol. 46, no. 11, pp. 100-109, Nov. 2008.

Security and Privacy for VC Systems (cont'd)

Beacon packet

1. Generate signature with SK_1
2. Append certificate
3. Send packet



1. Validate certificate (if not previously done so)
2. Validate signature
3. Validate geo-stamp in the header
4. Accept/Reject packet



- ▶ Sign packets with the private key, corresponding to the current valid pseudonym
- ▶ Verify packets with the valid pseudonym
- ▶ Cryptographic operations in a Hardware Security Module (HSM)



Problem Statement and Motivation

The design of a VPKI

- ▶ Resilience
- ▶ Stronger adversarial model (than fully-trustworthy entities)
 - ▶ User privacy protection against *"honest-but-curious"* entities
 - ▶ User privacy enhancement and service unlinkability (inference of service provider or time)
- ▶ Pseudonym acquisition policies
 - ▶ How should each vehicle interact with the VPKI, e.g., how frequently and for how long?
 - ▶ Should each vehicle itself determine the pseudonym lifetime?
- ▶ Operation across multiple domains, thus a scalable design
- ▶ Efficiency and robustness



Security and Privacy Requirements for the VPKI Protocols

- ▶ Authentication, communication integrity and confidentiality
- ▶ Authorization and access control
- ▶ Non-repudiation, accountability and eviction (revocation)
- ▶ Privacy
 - ▶ Anonymity (conditional)
 - ▶ Unlinkability
- ▶ Thwarting Sybil-based misbehavior
- ▶ Availability



Adversarial Model

External adversaries

Internal adversaries

Stronger adversarial model

Protection against *honest-but-curious* VPKI entities

- ▶ Correct execution of protocols but motivated to profile users
- ▶ Concealing pseudonym provider identity and acquisition time, and reducing pseudonyms linkability (inference based on time)

Multiple VPKI entities could collude

Secure VC System

- ▶ Root Certification Authority (RCA)
- ▶ Long Term CA (LTCA)
- ▶ Pseudonym CA (PCA)
- ▶ Resolution Authority (RA)
- ▶ Lightweight Directory Access Protocol (LDAP)
- ▶ Roadside Unit (RSU)
- ▶ Trust established with RCA, or through cross certification

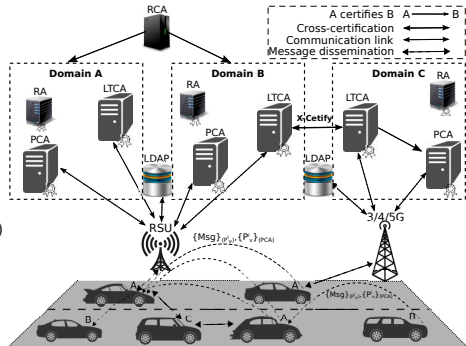


Figure: VPKI Overview

System Model

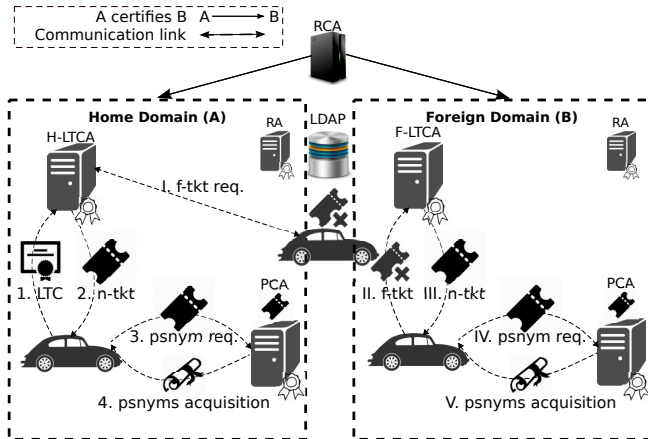


Figure: VPKI Architecture



Experimental Setup (#1)

▶ VPKI testbed

- ▶ Implementation in C++
- ▶ OpenSSL: Transport Layer Security (TLS) and Elliptic Curve Digital Signature Algorithm (ECDSA)-256 according to the standard

▶ Network connectivity

- ▶ Varies depending on the actual OBU-VPKI connectivity
- ▶ Reliable connectivity to the VPKI (e.g., RSU, Cellular, opportunistic WiFi)

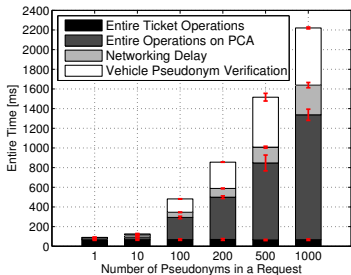
Table: Servers and Clients Specifications

	LTCA	PCA	RA	Clients
VM Number	2	5	1	25
Dual-core CPU (Ghz)	2.0	2.0	2.0	2.0
BogoMips	4000	4000	4000	4000
Memory	2GB	2GB	1GB	1GB
Database	MySQL	MySQL	MySQL	MySQL
Web Server	Apache	Apache	Apache	-
Load Balancer	Apache	Apache	-	-
Emulated Threads	-	-	-	400

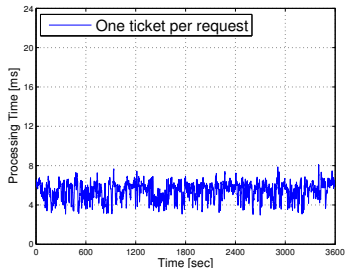
▶ Use cases

- ▶ Pseudonym provision
- ▶ Performing a DDoS attack

Client and LTCA Performance Evaluation



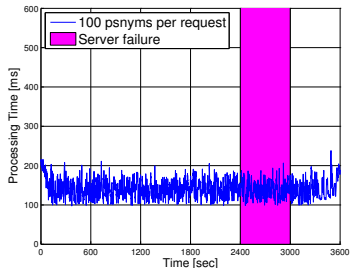
Client processing time



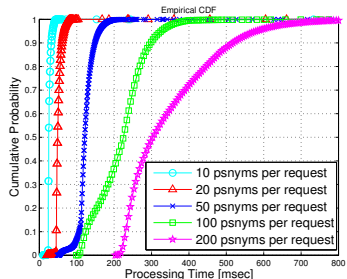
LTCA performance

- ▶ Delay to obtain pseudonyms
- ▶ LTCA response time to issue a ticket

PCA Performance Evaluation



Issuing 100 pseudonyms per request



PCA performance under different configuration

- ▶ PCA response time, including a *crash* failure
- ▶ Efficient provision for pseudonyms, with different configurations
- ▶ Obtaining 200 pseudonyms: $F_x(t=500)=0.9$ or $\Pr\{t \leq 500\}=0.9$



Experimental Setup (#2)

Table: Mobility Traces Information

	TAPASCologne	LuST
Number of vehicles	75,576	138,259
Number of trips	75,576	287,939
Duration of snapshot (hour)	24	24
Available duration of snapshot (hour)	2 (6-8 AM)	24
Average trip duration (sec.)	590.49	692.81
Total trip duration (sec.)	44,655,579	102,766,924

► Main metric

- End-to-end pseudonym acquisition latency from the initialization of ticket acquisition protocol till successful completion of pseudonym acquisition protocol

Table: Servers & Clients Specifications

	LTCA	PCA	Client
Number of entities	1	1	1
Dual-core CPU (Ghz)	2.0	2.0	2.0
BogoMips	4000	4000	4000
Memory	2GB	2GB	1GB
Database	MySQL	MySQL	MySQL

- N.B. PRESERVE Nexcom boxes specs:
dual-core 1.66 GHz, 2GB Memory

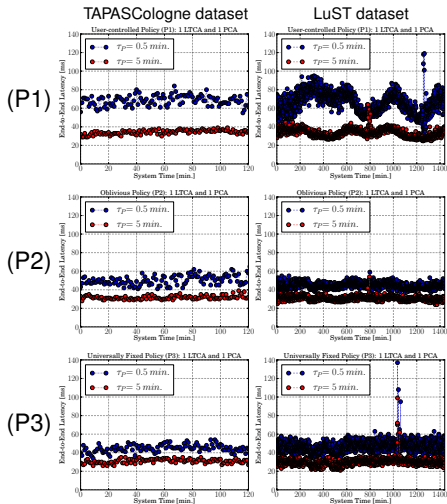
End-to-end Latency for P1, P2, and P3

Choice of parameters:

- Frequency of interaction and volume of workload to a PCA
- $\Gamma=5$ min., $\tau_P=0.5$ min., 5 min.

LuST dataset ($\tau_P = 0.5$ min):

- P1: $F_x(t = 167 \text{ ms}) = 0.99$
- P2: $F_x(t = 80 \text{ ms}) = 0.99$
- P3: $F_x(t = 74 \text{ ms}) = 0.99$





Summary of Contributions

1. Facilitating multi-domain operation
2. Offering increased user privacy protection
 - ▶ Honest-but-curious system entities
 - ▶ Eliminating pseudonym linking based on timing information
3. Eradication of Sybil-based misbehavior
4. Proposing multiple generally applicable pseudonym acquisition policies
5. Detailed analysis of security and privacy protocols
6. Extensive experimental evaluation
 - ▶ Efficiency, scalability, and robustness
 - ▶ Achieving significant performance improvement
 - ▶ Modest VMs can serve sizable areas or domain



Original Work

- ▶ N. Alexiou, M. Laganà, S. Gisdakis, M. Khodaei, and P. Papadimitratos, “*VeSPA: Vehicular Security and Privacy-preserving Architecture*,” in ACM HotWiSec, Budapest, Hungary, Apr. 2013.
- ▶ M. Khodaei, H. Jin, and P. Papadimitratos, “*Towards Deploying a Scalable & Robust Vehicular Identity and Credential Management Infrastructure*,” in IEEE VNC, Paderborn, Germany, Dec. 2014.
- ▶ M. Khodaei and P. Papadimitratos, “*The Key to Intelligent Transportation: Identity and Credential Management in Vehicular Communication Systems*,” IEEE VT Magazine, vol. 10, no. 4, pp. 63-69, Dec. 2015.
- ▶ M. Khodaei and P. Papadimitratos, “*Evaluating On-demand Pseudonym Acquisition Policies in Vehicular Communication Systems*,” in ACM MobiHoc, Workshop on Internet of Vehicles and Vehicles of Internet (IoV-Vol), Paderborn, Germany, July 2016.
- ▶ M. Khodaei, H. Jin, and P. Papadimitratos, “*SECMACE: Scalable and Robust Identity and Credential Management Infrastructure in Vehicular Communication Systems*,” Submitted to the IEEE Transactions on Intelligent Transportation Systems.



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ITRL Conference on Integrated Transport

Mohammad Khodaei
Networked Systems Security Group (NSS)
www.ee.kth.se/nss