

# Towards Deploying a Scalable & Robust Vehicular Identity and Credential Management Infrastructure

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# Outline

## 1 Introduction

- Background
- Contributions
- System Outline

## 2 Our Solution

- System Overview
- VPKI Services & Protocols

## 3 Security Analysis

## 4 Performance Evaluation

## 5 Conclusions



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# Background

## Projects

SEVECOM, EVITA, PRECIOSA, OVERSEE, DRIVE-C2X, PRESERVE,  
CAMP-VSC3

## Standardization and Harmonization

IEEE 1609.2, ETSI and C2C-CC: Vehicular Communication (VC) related specifications for privacy-preserving architectures

## Vehicular Public-Key Infrastructure (VPKI)

- Cornerstone for all these efforts
- Consensus on the need and basic characteristics



# Background (Cont'd)

## Inference of service provider or service time

- C2C-CC [1], PRESERVE [2], SCMS [3]
- CoPRA [4] linking pseudonyms to the real-identity
- VeSPA [5, 6], SEROSA [7]
- V-Token [8] learns the real identity of V-Token's owner

## Sybil-based misbehavior not precluded by VPKI

- C2C-CC [1], PRESERVE [2], SCMS [3]
- VeSPA [5, 6], SEROSA [7]
- **One remedy**
  - Non-overlapping pseudonym lifetimes
  - Downside: easier linkability



# Contributions

## Stronger adversarial model

Increased protection against *honest-but-curious* VPKI entities

- No inference of service provider or time
- Correct execution of protocols, but motivated to profile users

## Eradication of Sybil-based misbehavior

VPKI design that ensures a compromised vehicle cannot obtain multiple pseudonyms valid simultaneously

## Extensive & detailed experimental evaluations

- Full-blown VC *standard-compliant* implementation of VPKI
- Significant performance improvement
  - Multi-domain operation
  - Efficiency
  - Scalability

# System Outline

- Vehicles registered with one **Long Term Certification Authority (LTCA)** (home domain)
- **Pseudonym Certification Authority (PCA)** servers in one or multiple domains
- Vehicles can obtain pseudonyms from any **PCA** (in home or foreign domains)
- Trust with the help of a **Root Certification Authority (RCA)**
- Trust associations of PCAs and LTCAs through **Lightweight Directory Access Protocol (LDAP)** services
- Resolve a pseudonym with the help of a **Resolution Authority (RA)**



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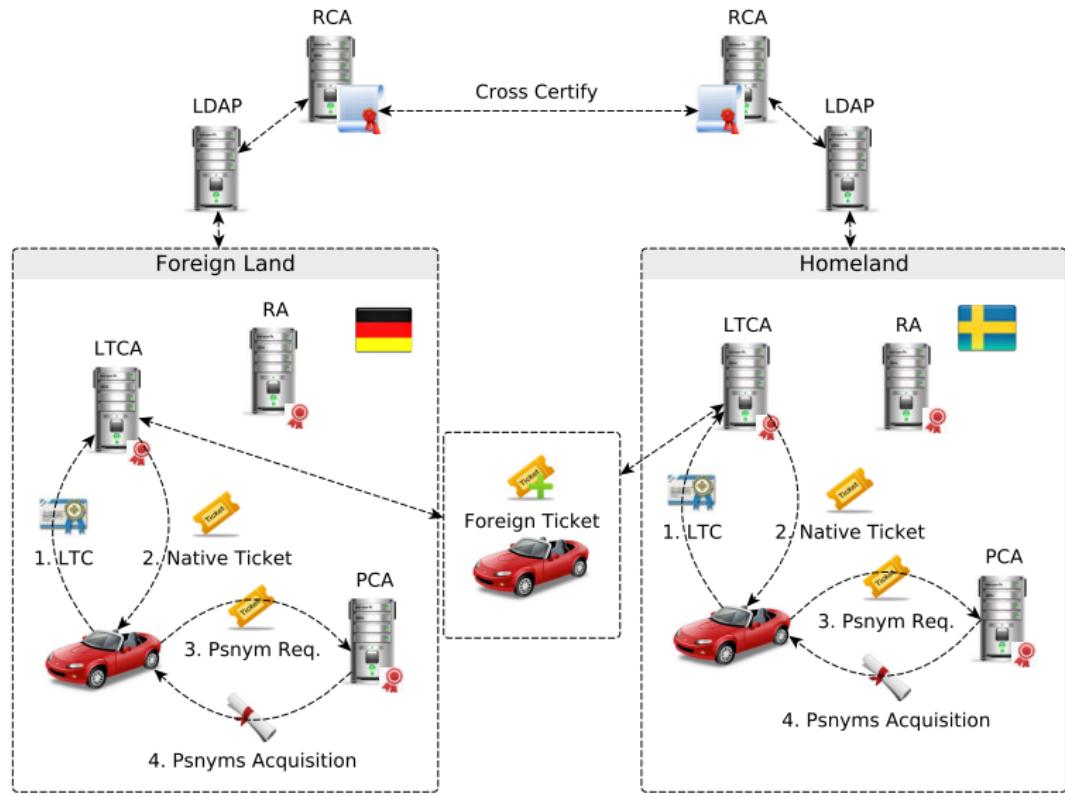
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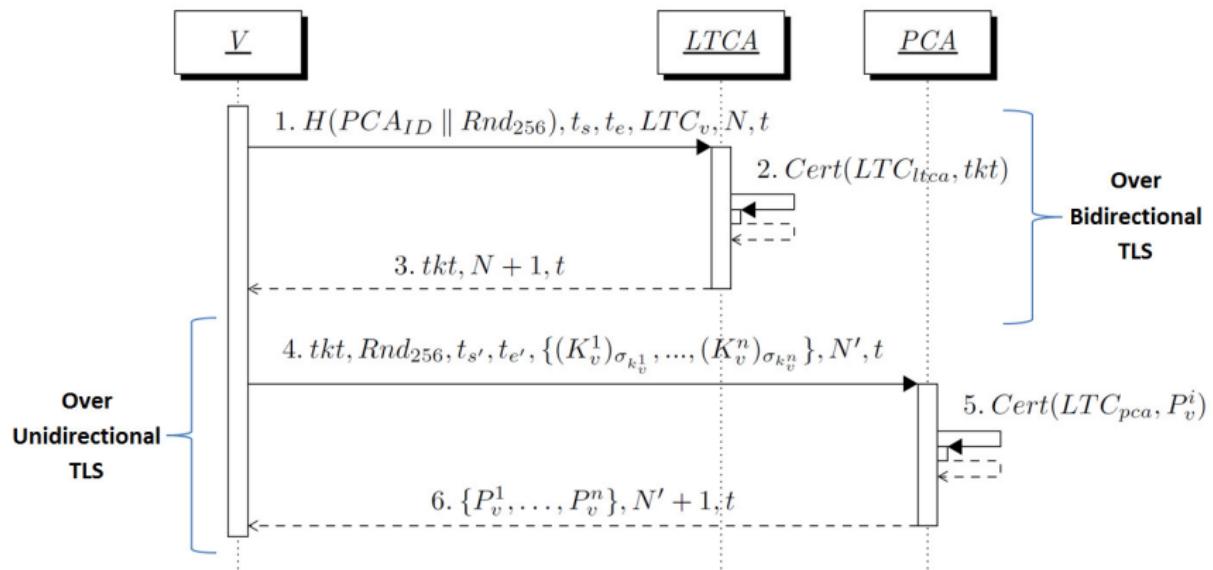
## 5 Conclusions



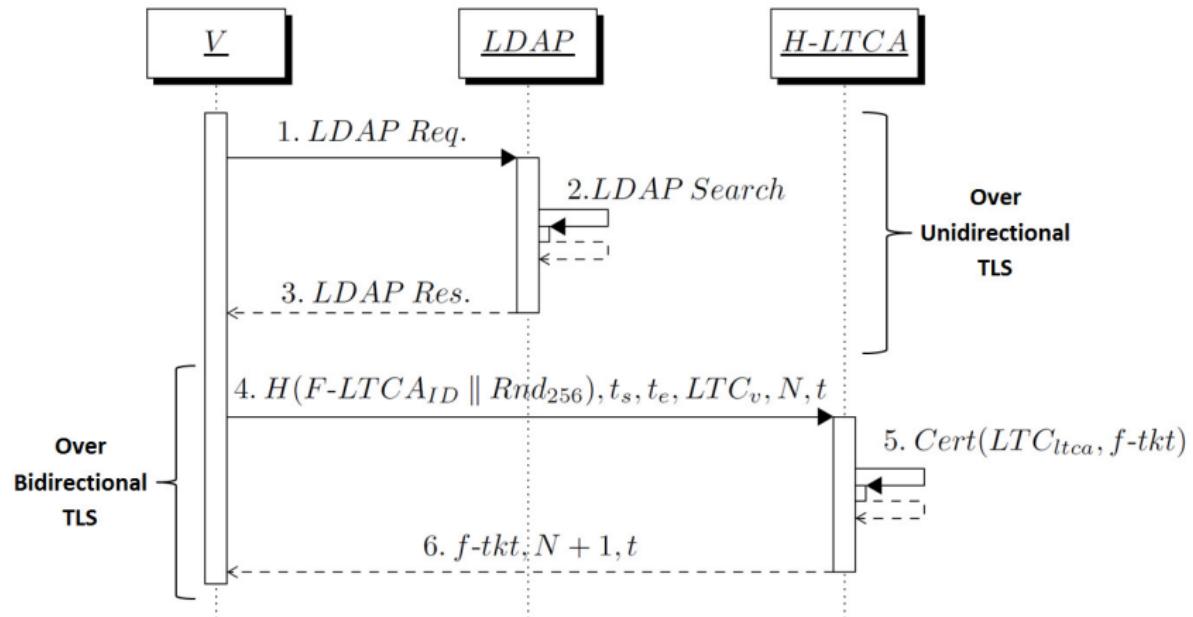
# VPKI Architecture



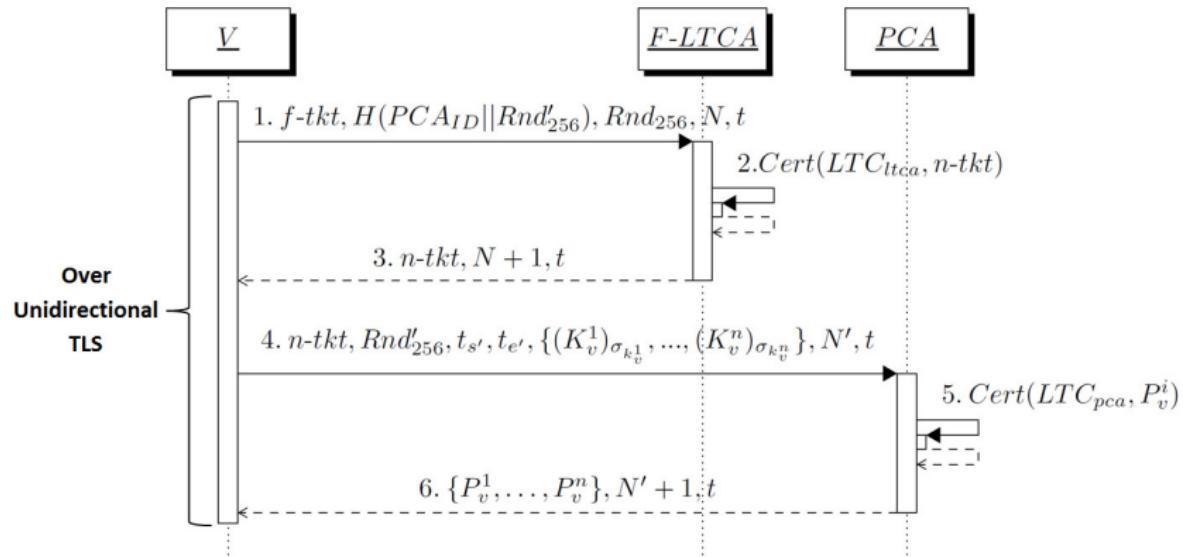
# Ticket and Pseudonym Acquisition



# Roaming User: Foreign Ticket Authentication



# Native Ticket and Pseudonym Acquisition in the Foreign Domain



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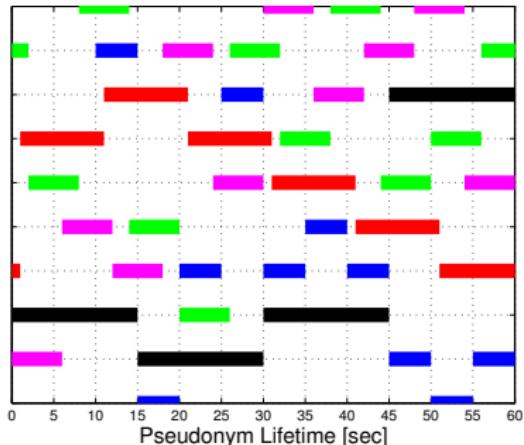
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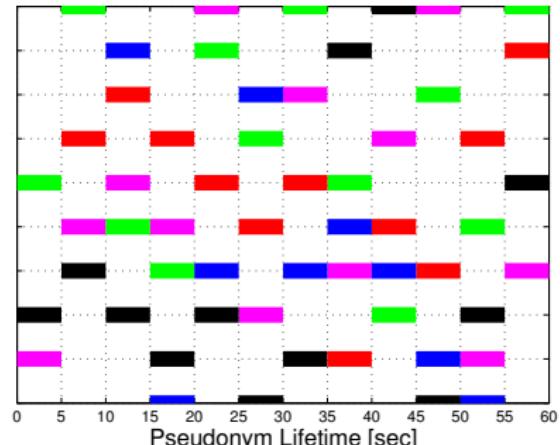
- **Communication integrity, confidentiality, non-repudiation**
  - Certificates, TLS and digital signatures
- **Authentication and authorization**
  - LTCA performs Authentication, Authorization and Accounting (AAA)
  - PCA grants the service
  - Security Association through LDAP
- **Concealing pseudonym providers, foreign identity providers and actual pseudonym acquisition period**
  - Sending  $H(PCA_{id} \parallel Rnd_{256})$ ,  $t_s$ ,  $t_e$ ,  $LTC_v$  to the home LTCA
  - PCA verifies if  $[t'_s, t'_e] \subseteq [t_s, t_e]$
- **Thwarting Sybil-based misbehavior**
  - LTCA keeps the records of the issued tickets
  - A ticket is bound to a specific PCA
  - PCA keeps records of ticket usage



# Ticket & Pseudonym Lifetime Policy



Flexible lifetimes



Fixed lifetimes

- Non-overlapping pseudonym lifetimes from eavesdroppers' perspective
- Distinct lifetimes per vehicle make linkability easier
- Uniform pseudonym lifetime in a domain
- No distinction among obtained pseudonyms set, thus less probable to link pseudonyms

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# Server and Client Specifications

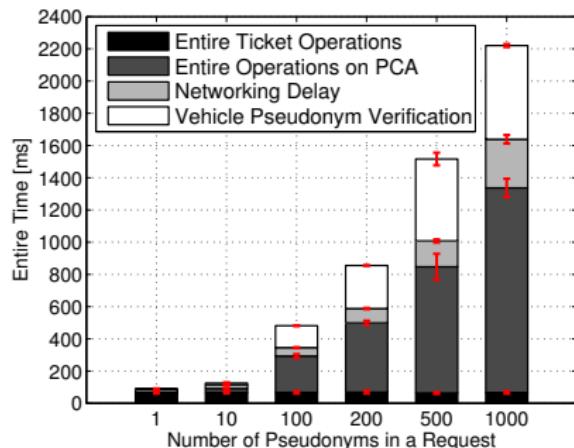
	LTCA	PCA	RA	Clients
VM Number	2	5	1	25
Dual-core CPU (Ghz)	2.0	2.0	2.0	2.0
BogoMips	4000	4000	4000	4000
Memory	2GB	2GB	1GB	1GB
Database	MySQL	MySQL	MySQL	MySQL
Web Server	Apache	Apache	Apache	-
Load Balancer	Apache	Apache	-	-
Emulated Threads	-	-	-	400

Use cases:

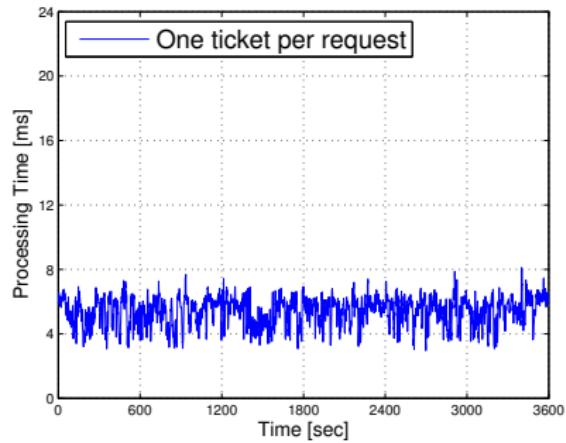
- Pseudonym provision
- Pseudonym resolution & revocation
- Performing DDoS attack



# Client and LTCA Performance Evaluation



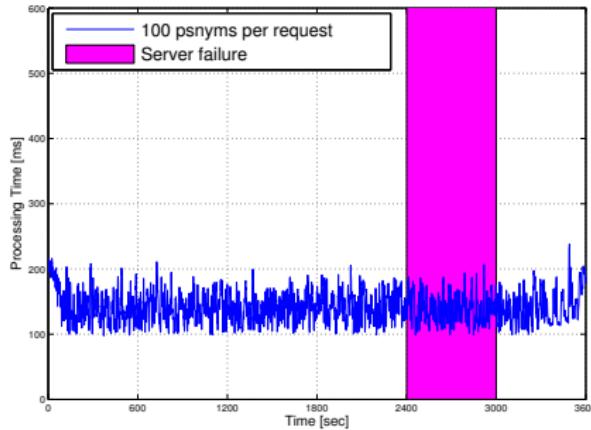
Client processing time



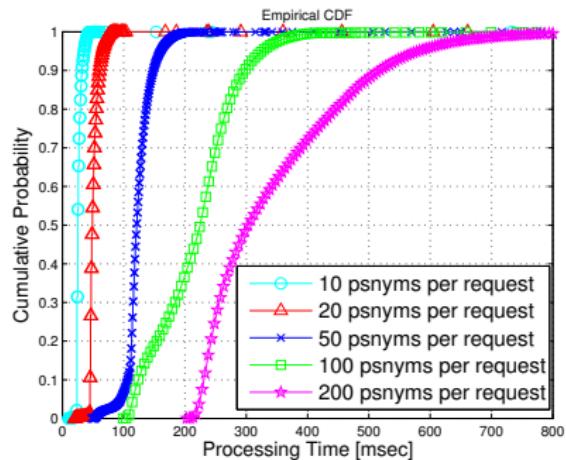
LTCA performance

- Delay to obtain pseudonyms
- LTCA response time to issue a ticket

# PCA Performance Evaluation



Issuing 100 pseudonyms per request

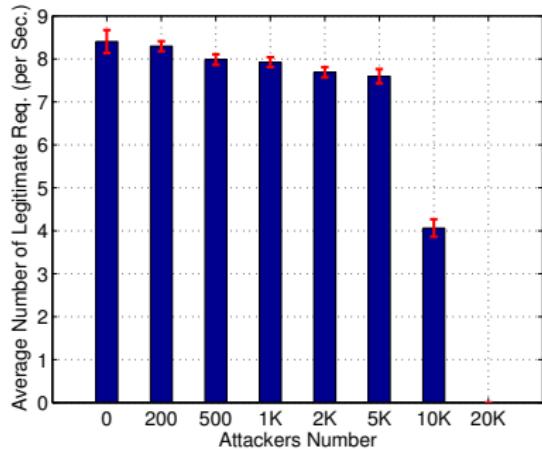


PCA performance under different configuration

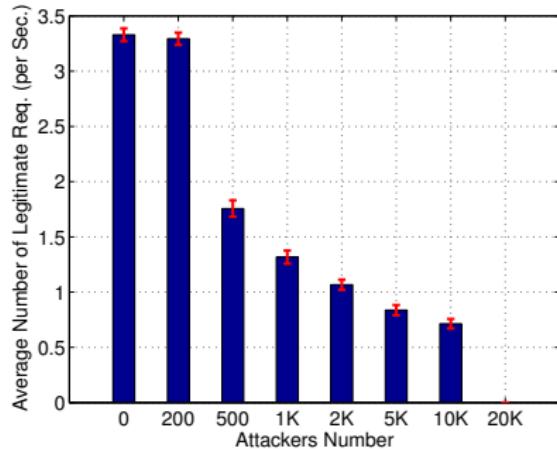
- PCA response time, including a *crash* failure
- Efficient provision for pseudonyms, with different configurations



# VPKI Servers under DDoS Attack



LTCA performance

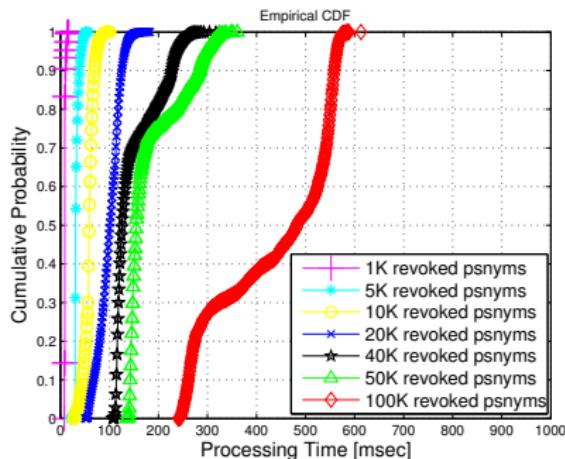


PCA performance

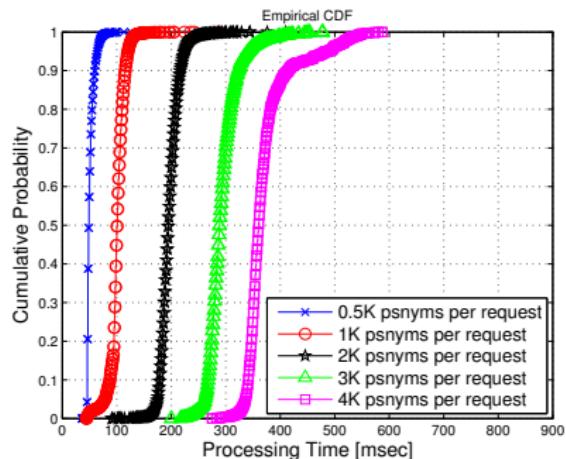
- 10K legitimate vehicles, requesting 100 pseudonyms every 10 minutes
- Up to 20K attackers, sending requests every 10 seconds
- An LTCA is more resistant to DDoS than a PCA



# Performance Evaluation for Pseudonym Revocation (CRL<sup>1</sup> or OCSP<sup>2</sup>) and Resolution



Obtaining a CRL



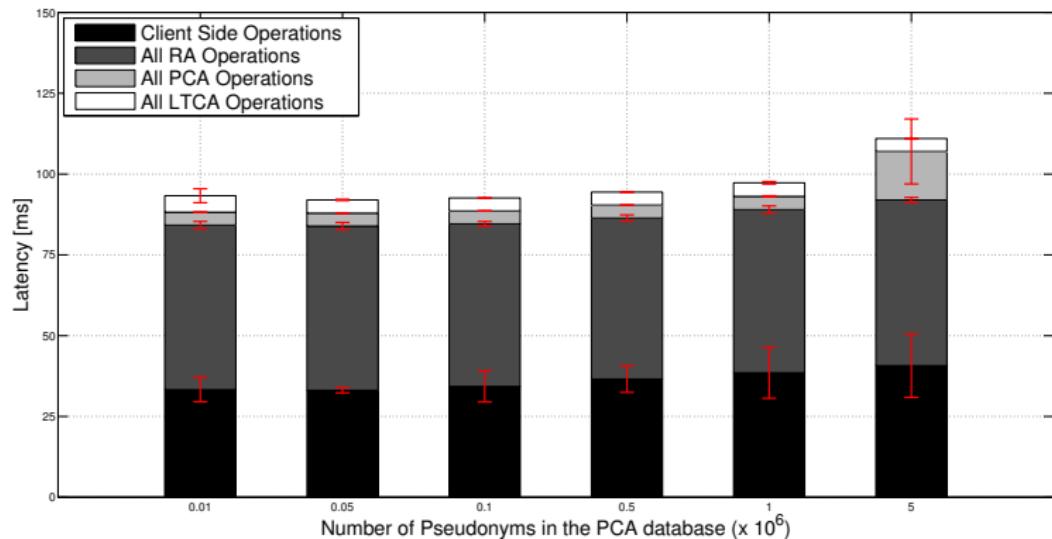
OCSP validation

- For 50K CRL:  $F_x(t=280)=0.9$  or  $\Pr\{t \leq 280\}=0.9$
- For 4K OCSP:  $F_x(t=400)=0.9$  or  $\Pr\{t \leq 400\}=0.9$

<sup>1</sup>CRL: Certificate Revocation List

<sup>2</sup>OCSP: Online Certificate Status Protocol

# Entities Response Time to Resolve & Revoke a Pseudonym



- On average 100 ms to resolve & revoke a pseudonym



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# Conclusion

- Strong adversarial model: *honest-but-curious* VPKI entities
- Prevention from Sybil-based misbehavior
- Reduction of pseudonym linkability, thus message linkability
- Extensive evaluation of a full-blown VC *standard compliant* VPKI
  - Very significant performance improvement over prior systems



# Bibliography

- [1] **Car-to-Car Communication Consortium (C2C-CC).** [Online]. Available: <http://www.car-2-car.org/>
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# Questions and Discussion

# Thank you!

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