



Computing relative Betti diagrams using Koszul complexes

Isaac Ren with W. Chachólski, A. Guidolin, M. Scolamiero, and F. Tombari
January 5, 2024 — JMM 2024

Summary

- In multiparameter persistence, we do not have a simple decomposition theorem like **barcode decompositions**.
- Instead, we approximate persistence modules by simpler modules using **relative projective resolutions**.
- Under certain conditions, we explicitly compute the **Betti diagrams** of these resolutions using **Koszul complexes**.

Computing relative Betti diagrams using Koszul complexes

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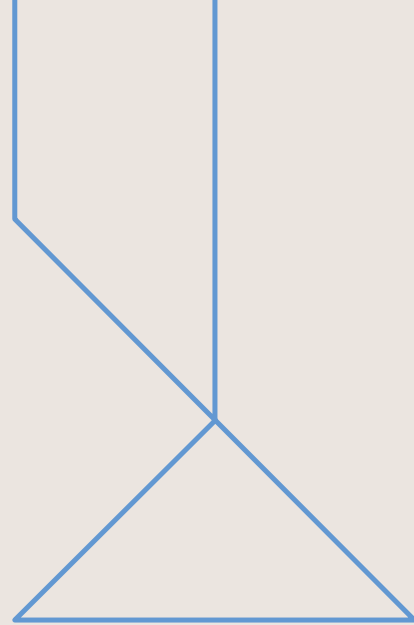
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Betti diagrams





Persistence modules as functors

- We consider functors $M: I \rightarrow \text{vect}_k$ where (I, \leq) is an arbitrary poset.
- We denote by $\text{Fun}(I, \text{vect}_k)$ the category of **functors indexed by I** .

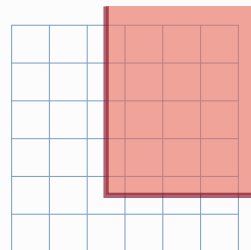
Persistence modules as functors

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- We denote by $\text{Fun}(I, \text{vect}_k)$ the category of **functors indexed by I** .
- For a in I , the **free functor at a** is the functor $k[a, \infty): I \rightarrow \text{vect}_k$ such that

$$k[a, \infty)(b) = \begin{cases} k & \text{if } b \geq a, \\ 0 & \text{otherwise,} \end{cases}$$

with identity transition maps.

- For example, for $I = \mathbf{N}^2$, the free functor at $(3, 2)$ is



Free resolutions

- A **free resolution** of a functor $M: I \rightarrow \text{vect}_{\mathbf{k}}$ is an exact sequence

$$\cdots \longrightarrow F_1 \longrightarrow F_0 \longrightarrow M \longrightarrow 0$$

where, for all $d \geq 0$, $F_d = \bigoplus_{a \in I} \mathbf{k}[a, \infty)^{\beta^d(a)}$.

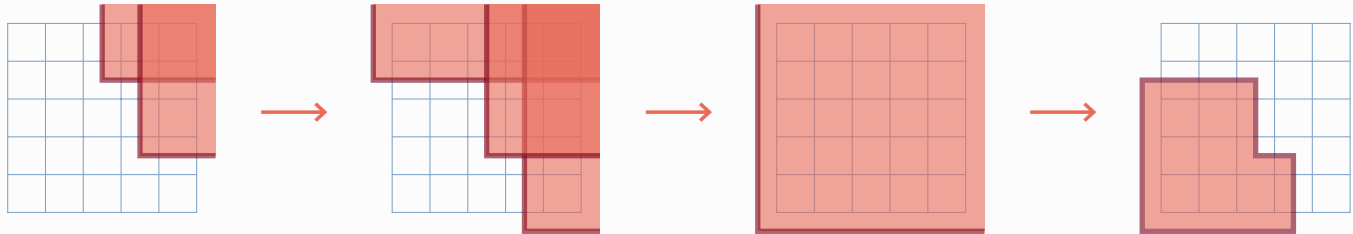
Free resolutions

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- Example:**



Betti diagrams

- We are interested in the unique **minimal free resolution**:

$$\dots \longrightarrow \bigoplus_{b \in I} \mathbf{k}[b, \infty)^{\beta^1(b)} \longrightarrow \bigoplus_{a \in I} \mathbf{k}[a, \infty)^{\beta^0(a)} \longrightarrow M \longrightarrow 0$$

Betti diagrams

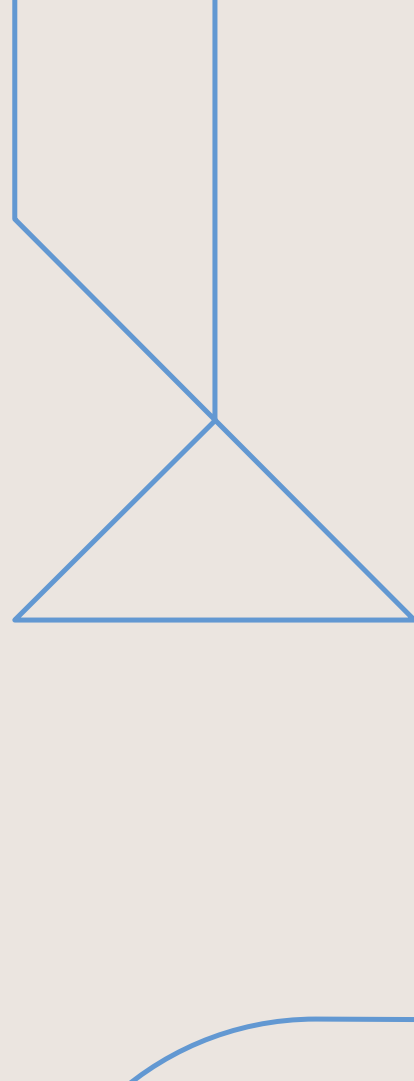
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- For all $d \geq 0$, we collect these multiplicities in a function $\beta^d M: I \rightarrow \mathbf{N}$ called the d^{th} **Betti diagram of M** .
- Problem:** In general, Betti diagrams require computing the entire minimal resolution. In particular, the differential maps are hard to compute.



Koszul complexes



Koszul complexes

- Suppose that (I, \leq) is an **upper semilattice** (all joins exist).
- For $M: I \rightarrow \text{vect}_k$ a functor and a in I , we define the **Koszul complex of M at a** as the chain complex $\mathcal{K}_a M$

$$\dots \longrightarrow \bigoplus_{\substack{b, c \text{ covers of } a \\ b \vee c \text{ exists}}} M(b \wedge c) \longrightarrow \bigoplus_{\substack{b \text{ cover of } a}} M(b) \longrightarrow M(a).$$

Koszul complexes

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- More formally, for all $d \geq 0$,

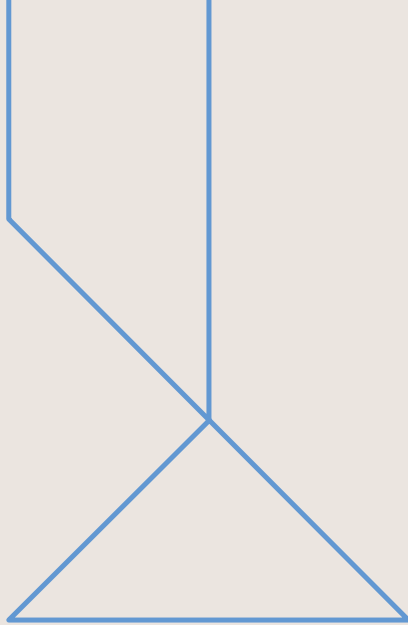
$$(\mathcal{K}_a M)_d := \bigoplus_{\substack{S \text{ subset of covers of } a \\ |S|=d \\ S \text{ has lower bound}}} M(\bigwedge_{(I \leq a)} S),$$

and the differential maps of $\mathcal{K}_a M$ are induced from the transition maps of M .

Theorem [Chachólski-Jin-Tombari 2021]

- Let (I, \leq) be a finite upper semilattice.
- For all functors $M: I \rightarrow \text{vect}_k$, elements a in I , and $d \geq 0$,

$$\beta^d M(a) = \dim H_d(\mathcal{K}_a M).$$

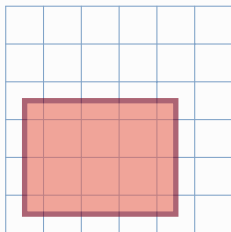


Relative homological algebra

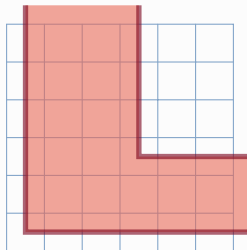


Relative resolutions

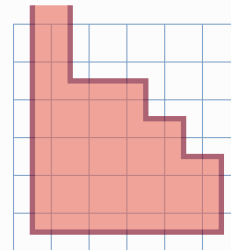
- Instead of resolving with free functors, we can try recreating bars.
- When $I = \mathbf{N}^2$, we can try:



rectangles



lower hooks
[BOO2022]



single-source spread
modules
[BBH2022]

Parameterization

- Let (J, \leq) be a poset.
- Let $\mathcal{T}: J^{\text{op}} \rightarrow \text{Fun}(I, \text{vect}_k)$ be a **parameterization** functor associating to each element a of J a functor $\mathcal{T}(a): I \rightarrow \text{vect}_k$.
- The collection of functors of interest is $\mathcal{P} := \{\mathcal{T}(a) \mid a \in J, \mathcal{T}(a) \neq 0\}$.
- \mathcal{T} is **thin** if, for all a, b in J , $\dim \text{Nat}(\mathcal{T}(a), \mathcal{T}(b)) \leq 1$.

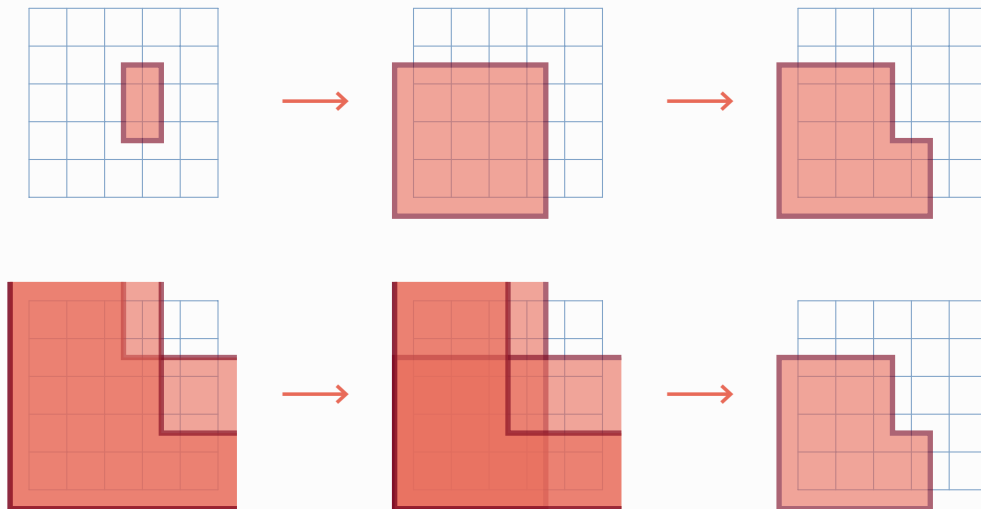
Relative projective resolutions

- Suppose that $\mathcal{T} : J^{\text{op}} \rightarrow \text{Fun}(I, \text{vect}_{\mathbf{k}})$ is thin.
- **\mathcal{P} -projective resolutions** are **\mathcal{P} -exact** sequences of direct sums of elements of \mathcal{P} :

$$\cdots \longrightarrow \bigoplus_{b \in J} \mathcal{T}(b)^{\beta_{\mathcal{P}}^1(b)} \longrightarrow \bigoplus_{a \in J} \mathcal{T}(a)^{\beta_{\mathcal{P}}^0(a)} \longrightarrow M \longrightarrow 0.$$

- **Example:** with $J = I$ and $\mathcal{T}(a) = \mathbf{k}[a, \infty)$, we recover standard free resolutions.

Examples



Relative Betti diagrams and Koszul complexes

- Similarly to the standard case, we collect the multiplicities of elements of \mathcal{P} in the minimal \mathcal{P} -projective resolution in **\mathcal{P} -Betti diagrams** $\beta_{\mathcal{P}}^d M: J \rightarrow \mathbf{N}$.
- **Problem:** we want to compute the \mathcal{P} -Betti diagrams of a functor $M: I \rightarrow \text{vect}_{\mathbf{k}}$.
- **Solution:** we compute the standard Betti diagrams of the functor

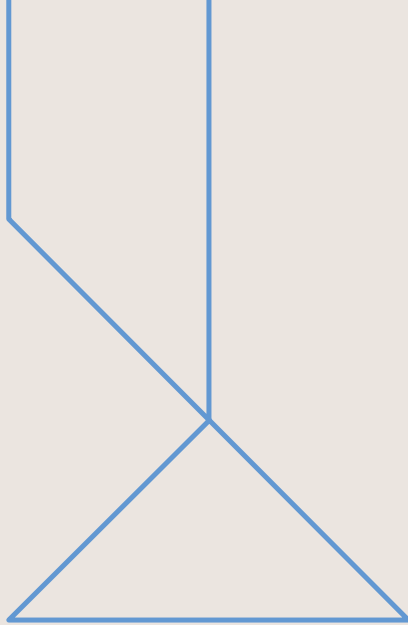
$$\text{Nat}(\mathcal{T}(-), M): \begin{cases} J \rightarrow \text{vect}_{\mathbf{k}} \\ a \mapsto \text{Nat}(\mathcal{T}(a), M) \end{cases}$$

using Koszul complexes, and then transfer the diagrams to the relative side.

Theorem [Chachólski, Guidolin, R., Scolamiero, Tombari]

- Let (J, \preceq) be a finite upper semilattice and \mathcal{T} a thin parameterization.
- Assume some other technical but realizable conditions.
- Then, for all functors $M: I \rightarrow \text{vect}_{\mathbf{k}}$, all a in J such that $\mathcal{T}(a) \neq 0$, and all $d \geq 0$,

$$\beta_{\mathcal{P}}^d M(a) = \dim H_d(\mathcal{K}_a \text{Nat}(\mathcal{T}(-), M)).$$



Computing an example



Lower hooks

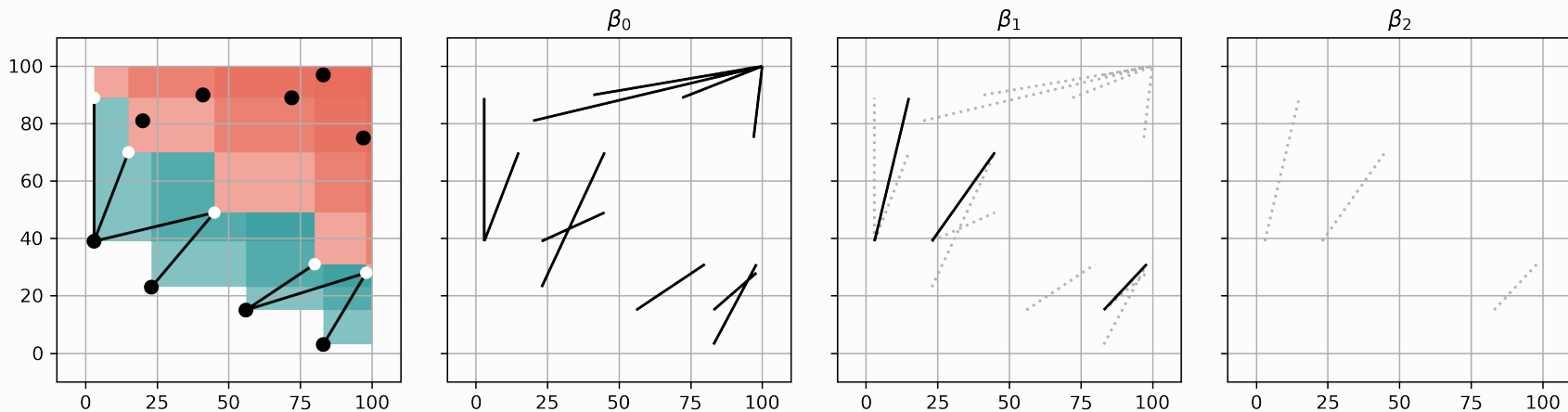
- We use the poset $J := \{(u, v) \mid u \leq v \in I\}$ with the product order, and the parameterization $\mathcal{T}: (u, v) \mapsto \text{coker}(\mathbf{k}[v, \infty) \rightarrow \mathbf{k}[u, \infty))$:

$$\mathcal{T}(u, v) =$$

- Fact:** For all (u, v) , we have $\text{Nat}(\mathcal{T}(u, v), M) = \ker(M(u) \rightarrow M(v))$.

Implementation

- We take as input a free presentation of a persistence module in the field of two elements \mathbf{F}_2 .



Complexity

- Time complexity is $O(2^{6r} n^{2r+3})$, where r is the dimension of the grid and n is the number of generators and relations in the input presentation.
- Space complexity is $O(rn^{2r+2})$.
- To compare, naively computing the free minimal resolution of $\text{Nat}(\mathcal{T}(-), M)$ is $O(r^2 2^{2r} n^{8r})$ and computing the signed barcode [BOO2022] is $O(r^3 n^{2r+1})$.

Cokernels vs. kernels

Proposition (Guidolin, R.)

- Let $M: I \rightarrow \text{vect}_k$ be a functor.
- Consider the induced functors $\ker M := (\ker(M(u) \rightarrow (v)))_{(u,v) \in J}$ and $\text{coker } M := (\text{coker}(M(u) \rightarrow (v)))_{(u,v) \in J}$ on J .
- Then, for all (u, v) in J ,

$$H_*(\mathcal{K}_{u,v}(\ker M)) \cong H_{*+2}(\mathcal{K}_{u,v}(\text{coker } M)).$$

Cokernels vs. kernels

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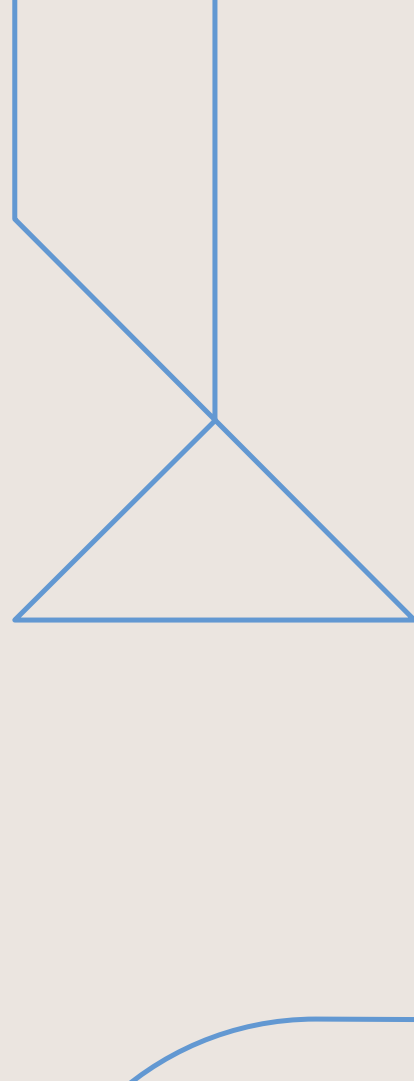
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- Then, for all (u, v) in J ,

$$H_*(\mathcal{K}_{u,v}(\ker M)) \cong H_{*+2}(\mathcal{K}_{u,v}(\text{coker } M)).$$

- Time complexity stays the same but space complexity goes down to $O(n^{2r+1})$.



Conclusions





Future work

- **Stability** and **hierarchical stabilization** of relative Betti diagrams.
- Construction of new **computable metrics** for functors.
- Relative resolutions over **continuous** posets.
- Comparing with relative **injective** resolutions



Thank you for your attention :)

References

- H. Asashiba, E. G. Escobar, K. Nakashima, and M. Yoshiwaki. *Approximation by interval-decomposables and interval resolutions of persistence modules*, 2023.
- B. Blanchette, T. Brüstle, and E. Hanson. *Homological approximations in persistence theory*, 2022.
- M. Botnan, S. Oppermann, and S. Oudot. *Signed barcodes for multi-parameter persistence via rank decompositions and rank-exact resolutions*, 2022.
- **Preprint:** *Koszul complexes and relative homological algebra of functors over posets*, arXiv:2209.05923.
- **Implementation:** <https://github.com/th-rtyf-re/lowerhooks>