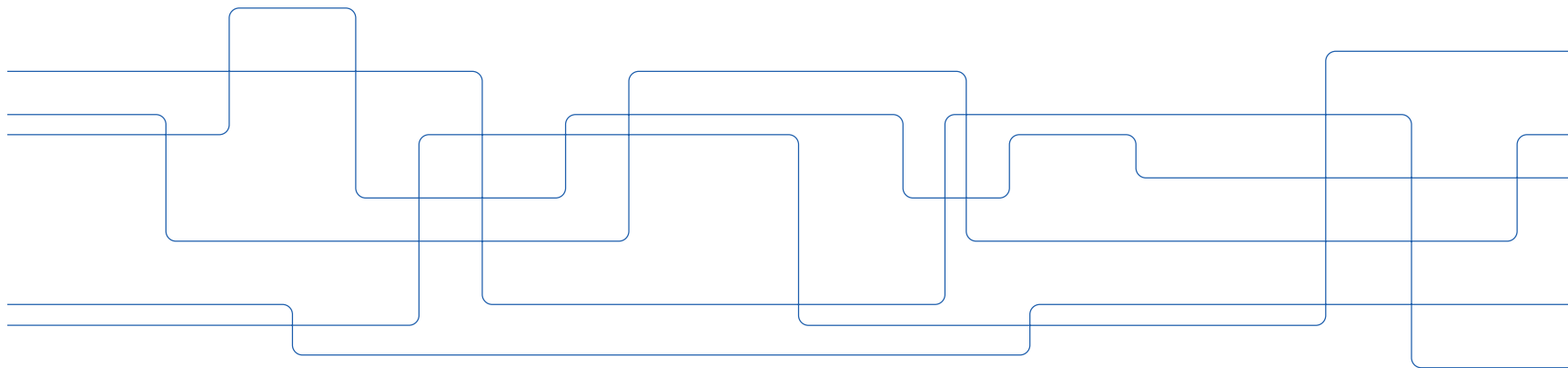


Relative homological algebra and Koszul complexes for multiparameter persistence

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Motivation

- › Consider a topological space X and n continuous real-valued functions $f_i: X \rightarrow \mathbf{R}, i \in \{1, \dots, n\}$.
- › For all a in \mathbf{R}^n , define $X_a := \{x \in X \mid \forall i \in \{1, \dots, n\}, f_i(x) \leq a_i\}$.
- › For all $d \geq 0$, we can study the d^{th} homology of the X_a 's.
 - › Moreover, if $a \leq b$ in \mathbf{R}^n for the product order, then the containment $X_a \subseteq X_b$ induces a linear map $H_d(X_a) \rightarrow H_d(X_b)$.
- › **Question:** What simple invariants can we compute from $H_d(X_\bullet): \mathbf{R}^n \rightarrow \mathbf{vect}_k$?

Today's talk

- › We can approximate persistence modules by simpler modules using **relative projective resolutions**.
- › Under certain conditions, we can explicitly compute the **Betti diagrams** of these resolutions using **Koszul complexes**.

Single-parameter persistence

Persistence modules as functors

- › Consider the poset (\mathbf{N}, \leq) .
- › We study persistence modules as **functors** $M: \mathbf{N} \rightarrow \mathbf{vect}_k$:
 - › For each natural number a in \mathbf{N} , we associate a k -vector space $M(a)$.
 - › For each pair $a \leq b$ in \mathbf{N} , we associate a linear map $M(a \leq b): M(a) \rightarrow M(b)$ called a **transition map**.
 - › For each triple $a \leq b \leq c$ in \mathbf{N} , $M(a \leq c) = M(b \leq c)M(a \leq b)$.

Natural transformations between functors

- › A **natural transformation** $f: M \rightarrow N$ between two functors M and N is
 - › the data, for each a in \mathbf{N} , of a linear map $f(a): M(a) \rightarrow N(a)$,
 - › such that, for each pair $a \leq b$ in \mathbf{N} , $f(b)M(a \leq b) = N(a \leq b)f(a)$:

$$\begin{array}{ccc}
 M(a) & \xrightarrow{M(a \leq b)} & M(b) \\
 f(a) \downarrow & & \downarrow f(b) \\
 N(a) & \xrightarrow{N(a \leq b)} & N(b)
 \end{array}$$

- › We denote by $\text{Nat}(M, N)$ the space of natural transformations from M to N .

Free functors

› For a in \mathbf{N} , the **free functor at a** is the functor $\mathbf{k}[a, \infty): \mathbf{N} \rightarrow \mathbf{vect}_{\mathbf{k}}$ such that

$$\mathbf{k}[a, \infty)(b) = \begin{cases} \mathbf{k} & \text{if } b \geq a, \\ 0 & \text{otherwise,} \end{cases}$$

with identity transition maps:

$$0 \rightarrow 0 \rightarrow \dots \rightarrow \underset{a-1}{0} \rightarrow \underset{a}{\mathbf{k}} \xrightarrow{\text{id}} \underset{a+1}{\mathbf{k}} \xrightarrow{\text{id}} \dots \xrightarrow{\text{id}} \mathbf{k} \xrightarrow{\text{id}} \dots$$

Free presentations

- › Every functor M can be presented as the quotient of two **free functors**:

$$0 \longrightarrow \bigoplus_{b \in \mathbb{N}} \mathbf{k}[b, \infty)^{\beta^1(b)} \longrightarrow \bigoplus_{a \in \mathbb{N}} \mathbf{k}[a, \infty)^{\beta^0(a)} \longrightarrow M \longrightarrow 0.$$

- › This sequence is **minimal** if all of its endomorphisms are isomorphisms.
- › **Fact:** the minimal presentation is unique.

Barcodes

We can do better than free presentations:

For $a < b$ in \mathbf{N} , the **bar from a to b** is the functor $\mathbf{k}[a, b]: \mathbf{N} \rightarrow \mathbf{vect}_{\mathbf{k}}$ such that

$$\mathbf{k}[a, b](c) = \begin{cases} \mathbf{k} & \text{if } a \leq c < b, \\ 0 & \text{otherwise,} \end{cases}$$

with identity transition maps:

$$0 \rightarrow 0 \rightarrow \dots \rightarrow \underset{a-1}{0} \rightarrow \underset{a}{\mathbf{k}} \xrightarrow{\text{id}} \underset{a+1}{\mathbf{k}} \xrightarrow{\text{id}} \dots \xrightarrow{\text{id}} \underset{b-1}{\mathbf{k}} \rightarrow \underset{b}{0} \rightarrow \dots$$

Theorem [Zomorodian-Carlsson 2005]

› Every functor in $\text{Fun}(\mathbf{N}, \text{vect}_k)$ is isomorphic to a unique direct sum of bars.

Standard homological algebra for multiparameter persistence

Multiparameter functors

- › Instead of functors $\mathbf{N} \rightarrow \mathbf{vect}_k$, we can consider functors $\mathbf{N}^n \rightarrow \mathbf{vect}_k$.
- › The poset (\mathbf{N}^n, \leq) is equipped with the product order.

Functors over arbitrary posets

- › We now consider functors $M: I \rightarrow \mathbf{vect}_k$ where (I, \leq) is an arbitrary poset.
- › We denote by $\mathbf{Fun}(I, \mathbf{vect}_k)$ the category of **functors indexed by I** .

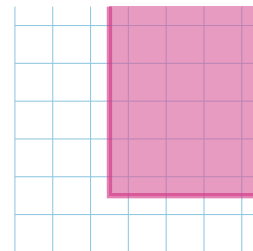
Free functors over posets

› For a in I , the **free functor at a** is the functor $\mathbf{k}[a, \infty): \mathbf{N} \rightarrow \mathbf{vect}_k$ such that

$$\mathbf{k}[a, \infty)(b) = \begin{cases} \mathbf{k} & \text{if } b \geq a, \\ 0 & \text{otherwise,} \end{cases}$$

with identity transition maps.

› For example, if $I = \mathbf{N}^2$, then the free functor at $(3, 2)$ is



Free resolutions

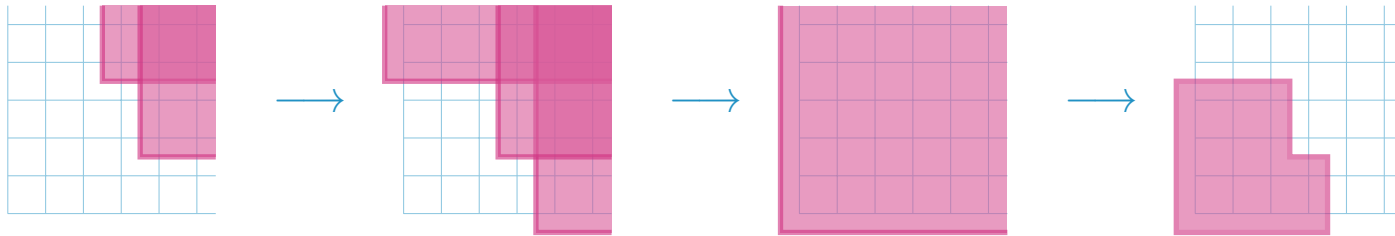
› A **free resolution** of a functor $M: I \rightarrow \mathbf{vect}_k$ is an exact sequence

$$\cdots \longrightarrow F_1 \longrightarrow F_0 \longrightarrow M \longrightarrow 0$$

where, for all $d \geq 0$, $F_d = \bigoplus_{a \in I} \mathbf{k}[a, \infty)^{\beta^d(a)}$.

› We also have the notion of a unique **minimal free resolution**.

Example



Betti diagrams

- › For a functor M , the multiplicities $\beta^d(a)$ of the unique minimal free resolution are of interest.
 - › For all $d \geq 0$, we collect these multiplicities in a function $\beta^d M: I \rightarrow \mathbf{N}$ called the d^{th} **Betti diagram of M** .
 - › **Problem:** In general, Betti diagrams require computing the entire minimal resolution. In particular, the differential maps are hard to compute.

Poset terminology

Let a and b be elements of the poset (I, \leq) .

- › A **cover of a** is a maximal element smaller than a .
- › The **join of a and b** , if it exists, is the unique minimal upper bound $a \vee b \geq a, b$.
- › The **meet of a and b** , if it exists, is the unique maximal lower bound $a \wedge b \leq a, b$.

Upper semilattices

- › The poset (I, \leq) is an **upper semilattice** if every pair of elements a, b has a join $a \vee b$.
- › If I is an upper semilattice, then we can individually compute values of Betti diagrams.

Koszul complexes for Betti diagrams

- › Suppose that (I, \leq) is an upper semilattice.
- › For $M: I \rightarrow \mathbf{vect}_k$ a functor and a in I , we define the **Koszul complex of M at a** as the chain complex $\mathcal{K}_a M$

$$\dots \longrightarrow \bigoplus_{\substack{b, c \text{ covers of } a \\ b \wedge c \text{ exists}}} M(b \wedge c) \longrightarrow \bigoplus_{b \text{ cover of } a} M(b) \longrightarrow M(a).$$

Standard homological algebra for multiparameter persistence

› More formally, for all $d \geq 0$,

$$(\mathcal{K}_a M)_d := \bigoplus_{\substack{S \text{ subset of covers of } a \\ |S|=d \\ S \text{ has lower bound}}} M(\bigwedge_{(I \leq a)} S).$$

› The differential maps of $\mathcal{K}_a M$ are induced from the transition maps of M .

Theorem [Chachólski-Jin-Tombari 2021]

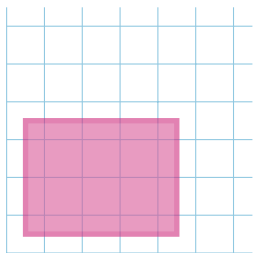
- › Let (I, \leq) be an upper semilattice.
- › For all functors $M: I \rightarrow \mathbf{vect}_k$, elements a in I , and $d \geq 0$,

$$\beta^d M(a) = \dim H_d(\mathcal{K}_a M).$$

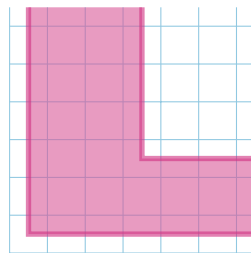
Relative homological algebra for multiparameter persistence

Non-free functors

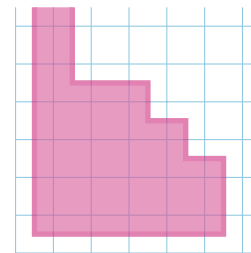
- › Instead of resolving with free functors, we can try recreating bars.
- › When $I = \mathbf{N}^2$, we can try:



rectangles



lower hooks
[BOO2021]



single-source spread
modules
[BBH2021]

Relative projectives

- › We fix a collection \mathcal{C} of functors in $\text{Fun}(I, \mathbf{vect}_k)$.
- › A natural transformation $f: M \rightarrow N$ is a **\mathcal{C} -epimorphism** if, for all A in \mathcal{C} , the linear map $\text{Nat}(A, f): \text{Nat}(A, M) \rightarrow \text{Nat}(A, N)$ is surjective.
- › A functor $A: I \rightarrow \mathbf{vect}_k$ is **\mathcal{C} -projective** if, for every \mathcal{C} -epimorphism $f: M \rightarrow N$, the linear map $\text{Nat}(A, f): \text{Nat}(A, M) \rightarrow \text{Nat}(A, N)$ is surjective.
- › A short sequence $L \rightarrow M \rightarrow N$ is **\mathcal{C} -exact** if, for all A in \mathcal{C} , the short sequence $\text{Nat}(A, L) \rightarrow \text{Nat}(A, M) \rightarrow \text{Nat}(A, N)$ is exact.

Relative projective resolutions

- › A **\mathcal{C} -projective resolution** of a functor $M: I \rightarrow \mathbf{vect}_k$ is a \mathcal{C} -exact sequence of functors

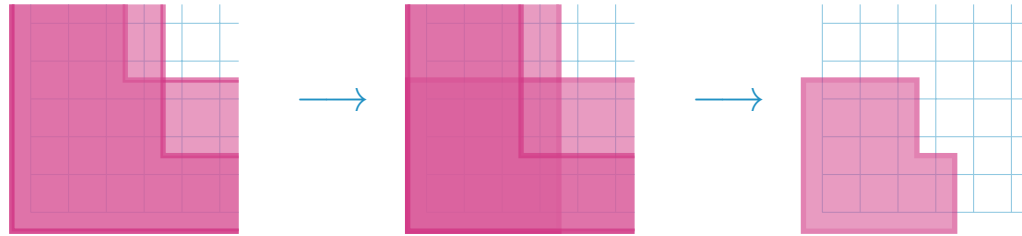
$$\cdots \longrightarrow F_1 \longrightarrow F_0 \longrightarrow M \longrightarrow 0$$

where the F_d are \mathcal{C} -projective.

- › We also have the notion of a unique **minimal \mathcal{C} -projective resolution**.

Example

› A lower hook resolution can look like



Parameterization by a poset

- › Let (J, \preceq) be a poset.
- › Let $\mathcal{P}: J^{\text{op}} \rightarrow \text{Fun}(I, \mathbf{vect}_k)$ be a functor associating to each element a of J a functor $\mathcal{P}(a): I \rightarrow \mathbf{vect}_k$.
 - › The collection of functors is now $\mathcal{C} := \{\mathcal{P}(a) \mid a \in J, \mathcal{P}(a) \neq 0\}$.
 - › \mathcal{P} is **thin** if, for all a, b in J , $\dim \text{Nat}(\mathcal{P}(a), \mathcal{P}(b)) \leq 1$.
 - › **Fact:** in this case, all \mathcal{C} -projectives are direct sums of elements of \mathcal{C} .

Relative Betti diagrams

- › Suppose that $\mathcal{P}: J^{\text{op}} \rightarrow \text{Fun}(I, \mathbf{vect}_k)$ is thin.
- › Relative projective resolutions are then sequences of direct sums of elements of \mathcal{C} :

$$\dots \longrightarrow \bigoplus_{b \in J} \mathcal{P}(b)^{\beta^1(b)} \longrightarrow \bigoplus_{a \in J} \mathcal{P}(a)^{\beta^0(a)} \longrightarrow M \longrightarrow 0.$$

- › Similarly to the standard case, we collect the multiplicities of elements of \mathcal{C} in the minimal \mathcal{C} -projective resolution in **\mathcal{P} -Betti diagrams** $\beta_{\mathcal{P}}^d M: J \rightarrow \mathbf{N}$.

Relative Koszul complexes

- › **Problem:** we want to compute the \mathcal{P} -Betti diagrams of a functor $M: I \rightarrow \mathbf{vect}_k$.
- › **Solution:** we compute the standard Betti diagrams of the functor

$$\mathrm{Nat}(\mathcal{P}(-), M): \begin{cases} J \rightarrow \mathbf{vect}_k \\ a \mapsto \mathrm{Nat}(\mathcal{P}(a), M) \end{cases}$$

using Koszul complexes, and then transfer the diagrams to the relative side.

Theorem

- › Let $M: I \rightarrow \mathbf{vect}_k$ be a functor and suppose that
 - › (J, \preceq) is an upper semilattice and \mathcal{P} is thin,
 - › $\{a \in J \mid \mathcal{P}(a) = 0\}$ is closed under joins,
 - › for all a, b in J , if a is minimal $\succeq b$ such that $\text{Nat}(\mathcal{P}(a), \mathcal{P}(b)) = 0$, then $\mathcal{P}(a) = 0$.
- › Then, for all a in J such that $\mathcal{P}(a) \neq 0$ and $d \geq 0$,

$$\beta_{\mathcal{P}}^d M(a) = \dim H_d(\mathcal{K}_a \text{Nat}(\mathcal{P}(-), M)).$$

Conclusions

Conclusions

Summary

Given a finite upper semilattice (J, \preceq) and a thin functor $\mathcal{P}: J^{\text{op}} \rightarrow \text{Fun}(I, \mathbf{vect}_k)$:

absolute poset

relative poset

$$(I, \leq)$$

$$(J, \preceq)$$

$$M: I \rightarrow \mathbf{vect}_k$$

$$\text{Nat}(\mathcal{P}(-), M): J \rightarrow \mathbf{vect}_k$$

copy of $\mathcal{P}(a)$ in the minimal
 \mathcal{C} -projective resolution

copy of $\mathbf{k}[a, -]$ in the minimal
free resolution

Koszul complexes of $\text{Nat}(\mathcal{P}(-), M)$ to compute multiplicities in
the minimal resolution

Outlook

- › **Software implementation** of the computation of Betti diagrams relative to lower hooks.
- › **Stability** and **hierarchical stabilization** of relative Betti diagrams.
- › Construction of new **computable metrics** for functors.

Thank you for your attention :)

References

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- › Preprint soon!