

# Complexity

Johan Montelius

KTH

HT23

# big-O

An estimate of the change in execution time...  
when the data set grows large.

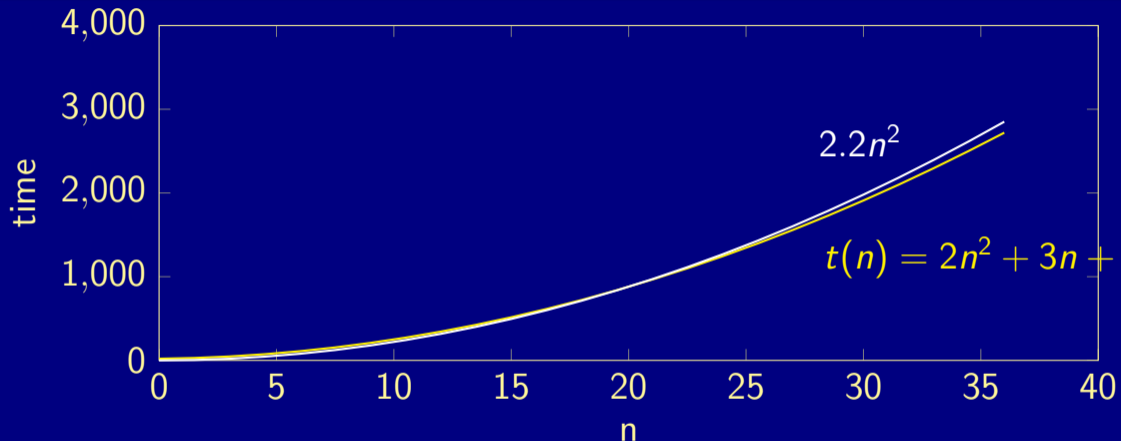
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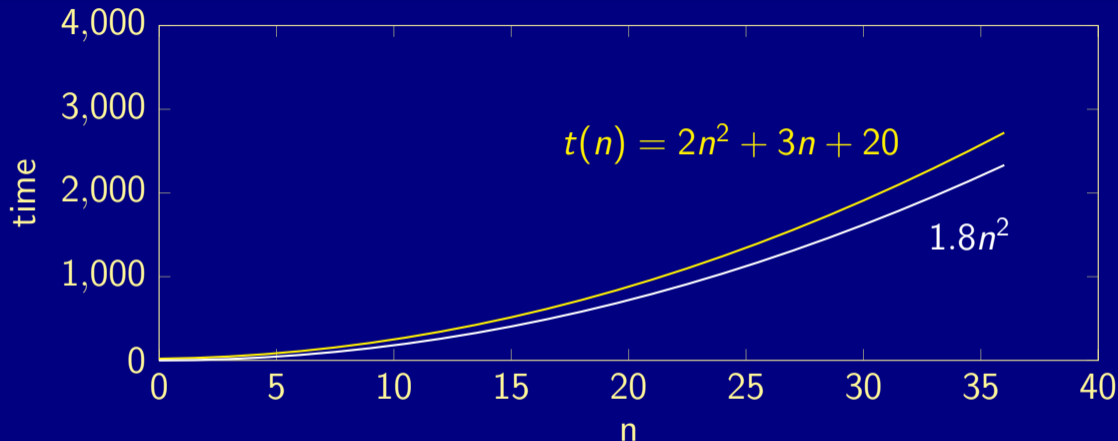
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# an upper bound



Ergo:  $t(n)$  is in  $O(n^2)$  since there is a  $k$  such that  $k \times n^2 > t(n)$  above some  $n$ .

# a lower bound



Ergo:  $t(n)$  is in  $\Omega(n^2)$  since there is a  $k$  such that  $k \times n^2 < t(n)$  above some  $n$ .

# Big-O, $\Omega$ and *Theta*

- A functions upper bound is limited by  $O(g(n))$ .
- Its lower bound is limited by  $\Omega(g(n))$ .
- If a function is limited by  $O(g(n))$  and  $\Omega(g(n))$  then it is limited by  $\Theta(g(n))$ .

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$$0.1 \times n + \log_{10}(n)$$

$$5 \times \sqrt{n} + 34$$

$$\ln(n)^2 + 34$$

$$\log_2(n) + 30$$

$$n^4 + 20 \times n + 32$$

$$50 \times n + 100$$

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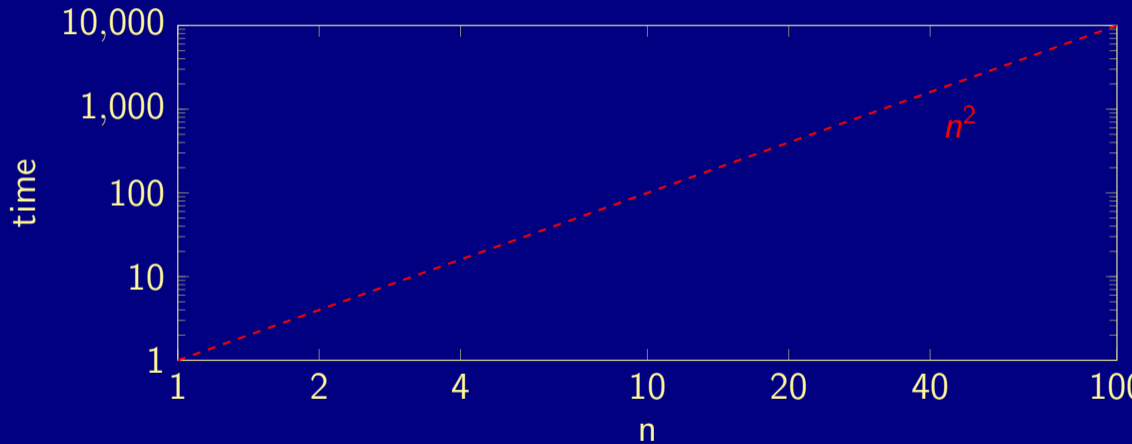
$$O(n^4) \quad O(n^2) \quad O(n \times \lg(n)) \quad O(n) \quad O(\sqrt{n}) \quad O(\lg(n)^2) \quad O(\lg(n))$$

orderd by complexity - not execution time given a specific  $n$

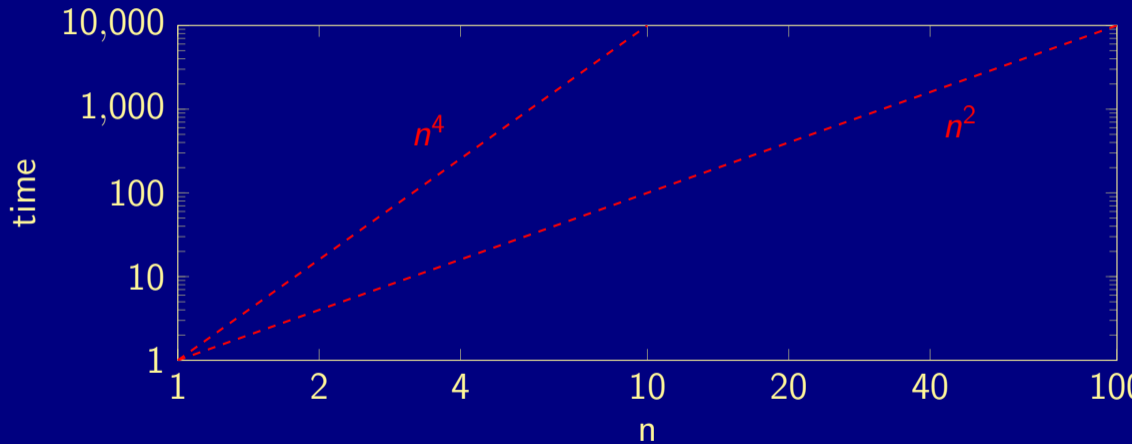
# the cost of chess



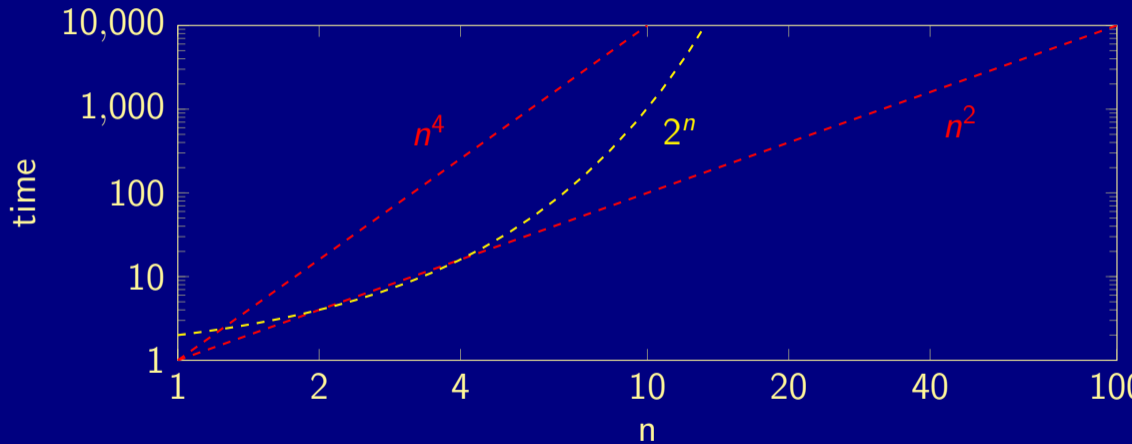
# exponential time



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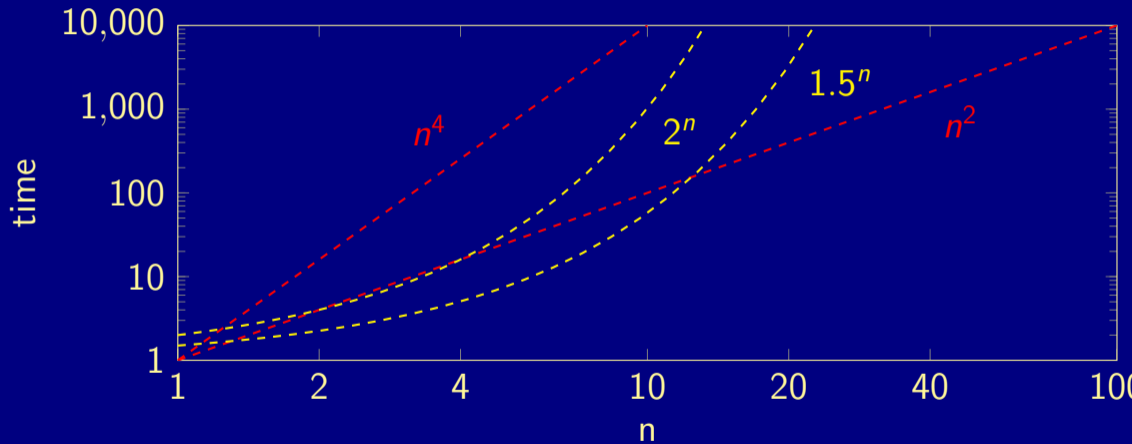


# exponential time

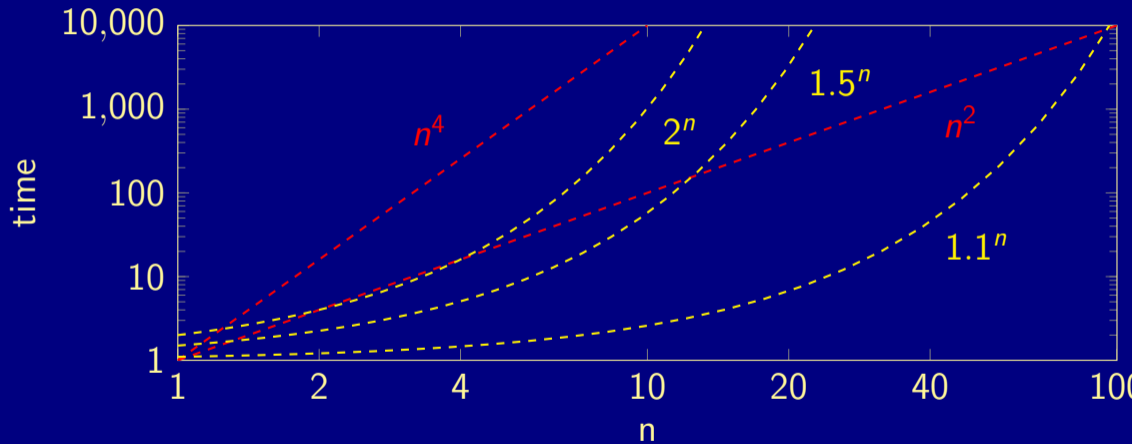




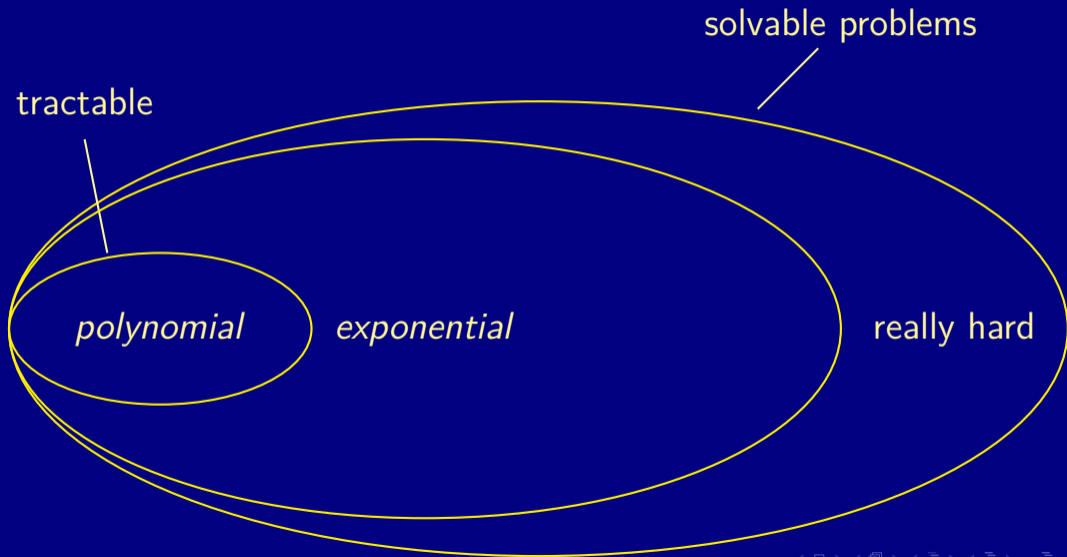
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# tractable, intractable and even worse



this sentence is false

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Kurt Gödel

*There are things  
that can not be  
decided.*

this sentence is false



Alonzo Church  
lambda calculus



Kurt Gödel  
*There are things  
that can not be  
decided.*

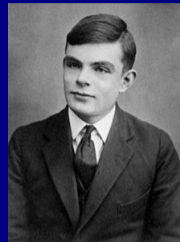
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Alan Turing  
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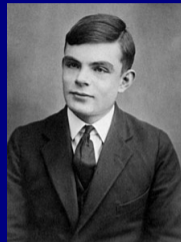
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Kurt Gödel  
*There are things  
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*Apart from things that can not be computed, we can compute  
everything :-)*



# The Imitation Game



# problem complexity

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If we have an algorithm at least we know the upper limit.

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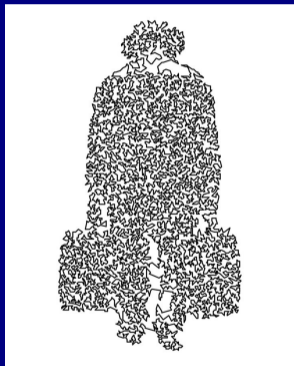
We have problems where it's easy to find an exponential algorithm, but no proof that there is no polynomial algorithm.

# traveling salesman

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- Find the shortest road that visits all cities.



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- Find the shortest road that visits all cities.
- Let's try them all ...

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- Find the shortest road that visits all cities.
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- If I give you the path - how can you verify that it is the shortest?

traveling salesman - almost the same

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- Find a road visiting all cities but that is less than  $x$  km.

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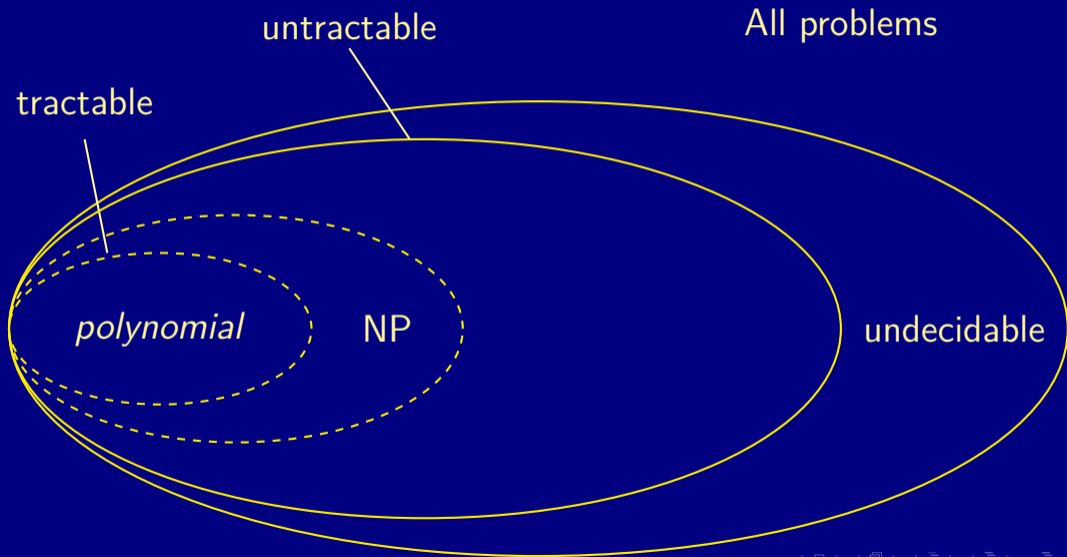
- Find a road visiting all cities but that is less than  $x$  km.
- Let's try them all ...
- .... exponential solution :-(  
● A solution is easy to verify in polynomial time.

# traveling salesman - almost the same



- Find a road visiting all cities but that is less than  $x$  km.
- Let's try them all ...
- .... exponential solution :-(  
• A solution is easy to verify in polynomial time.
- A Non-deterministic Polynomial problem - NP.

# non-deterministic polynomial



extra income

$$P = NP$$

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If you can prove it, or prove that it does not holds,

$$P = NP$$

If you can prove it, or prove that it does not hold,  
then you can claim a million dollar.

*The Millennium Prize Problems*

there is still hope



there is still hope

