FEL3330 Networked and Multi-agent Control Systems Lecture 10: Communication constraints: quantization, event-triggered sampling and control.

October 3, 2016

Today's lecture

- Quantized Consensus
- Event-triggered multi-agent control

Quantized consensus

- Each agent has quantized measurements of the form $q(x_i x_j)$, $q(\cdot)$ quantization function
- Can be extended to multiple dimensions
- Uniform and logarithmic quantizers

Quantizer types

- Uniform quantizer: $|q_u(a) a| \leq \delta_u, \forall a \in \mathbb{R}$
- Logarithmic quantizer: $\left|q_{l}\left(a\right)-a\right|\leq\delta_{l}\left|a\right|,orall a\in\mathbb{R}$



Closed-loop system

•
$$\dot{x}_i = u_i = -\sum_{j \in \mathcal{N}_i} a_{ij} q(x_i - x_j)$$

- Stack vector form: $\dot{\bar{x}} = -B^T B \Gamma q(\bar{x})$
- We use the positive definiteness of $B^T B$ when G is a tree graph

Theorem: Assume G is a tree.

- In the case of a uniform quantizer, the system converges to the set {x ∈ ℝ^N : |x_i − x_j| ≤ δ_a/2, {i,j} ∈ ε}.
- In the case of a logarithmic quantizer, the system asymptotically converges to x
 = 0, for all δ_l > 0.

Use $V = \bar{x}^T M^{-1} \bar{x}$, $M = B^T B$ as a candidate Lyapunov function

Quantized consensus

- Results extended to time varying communication topolgies (Guo and DVD, Automatica 2013)
- Additional properties of the (common) Lyapunov function $V = x^T H x$, $H = B(B^T B)^{-1} B^T$ are exploited

Event-triggered sampling

- Time-triggered sampling at pre-specified instants: does not take into account optimal resource usage
- A strategy considering better resource usage: event-triggered control
- Actuation updates in asynchronous manner
- Application to multi-agent systems

Event-triggered sampling

- Each agent *i* broadcast their state at discrete time instants t_0^i, t_1^i, \ldots
- $\hat{x}_i(t) = x_i(t_k^i)$: latest update for agent *i*
- Event-triggered control law: $u_i(t) = -\sum_{j \in N_i} (\hat{x}_i(t) \hat{x}_j(t))$
- Measurement error: $e_i(t) = \hat{x}_i(t) - x_i(t) = x_i(t_k^i) - x_i(t), t \in [t_k^i, t_{k+1}^i)$
- Next slides ack to Georg Seyboth

Event based scheduling of measurement broadcasts



Event-based broadcasting $\hat{x}_i(t) = x_i(t_k^i), t \in [t_k^i, t_{k+1}^i]$ $0 \le t_0^i \le t_1^i \le t_2^i \le \cdots$

- Consensus protocol $u_i(t) = -\sum_{i \in N_i} (\hat{x}_i(t) \hat{x}_j(t))$
- Measuremeent errors $e_i(t) = \hat{x}_i(t) x_i(t)$
- Closed-loop $\dot{x}(t) = -L\hat{x}(t) = -L(x(t) + e(t))$
- Disagreement $\delta(t) = x(t) a\mathbf{1}; \quad \mathbf{1}^T \delta(t) \equiv 0$

Trigger mechanism-Problem statement

• Define triger functions $f_i(\cdot)$ and triger when

 $f_i(t, x_i(t), \hat{x}_i(t), \hat{x}_{j_1}(t), \dots, \hat{x}_{j_{N_i}}(t)) > 0$

- Defines sequence of events $t_{k+1}^i = \inf\{t > t_k^i : f_i(t) > 0\}$
- Problem statement: Find suitable $f_i(\cdot)$ such that
 - no Zeno bevior occurs
 - · desired convergence properties are maintained
 - inter-agent communication is reduced as much as possible
- Intuition

$$f_i(e_i(t)) = |e_i(t)| - c_0 \Longrightarrow |e_i(t)| \le c_0, orall t \ge 0$$

Constant thresholds

$$\dot{x}(t) = u(t)$$
 $u(t) = -L\hat{x}(t)$ (1)

Theorem

Consider system (1) with undirected and connced graph G. Suppose that

1

$$f_i(e_i(t)) = |e_i(t)| - c_0$$

with $c_0 > 0$. Then, for all $x_0 \in \mathbb{R}^N$, the system does not exhibit Zeno behavior and for $t \to \infty$,

$$||\delta(t)|| \leq \frac{\lambda_N(G)}{\lambda_2(G)}\sqrt{N}c_0$$

Proof ideas:

Analytical solution of disagreement dynamics yields

$$||\delta(t)|| \leq e^{-\lambda_2(G)t}||\delta(0)|| + \lambda_N(G)\int_0^t e^{-\lambda_2(G)(t-s)}||e(s)||ds|$$

• Compute lower bound au on the inter-event intervals

Exponentially decreasing thresholds

Can we achieve asymptotic convergence?

Theorem

Consider system (1) with undirected and connected graph G. Suppose that

$$f_i(e_i(t)) = |e_i(t)| - c_1 e^{-at}$$

with $c_1 > 0$ and $0 < a < \lambda_2(G)$. Then, for all $x_0 \in \mathbb{R}^N$, the system does not exhibit Zeno behavior and for $t \to \infty$,

$$||\delta(t)||
ightarrow 0$$

Intuition:

λ₂(G) specifies the degree of convrgence for δ(t)

• $a < \lambda_2(G)$ implies that the threshold $c_1 e^{-at}$ decreases slower Remark: Vanishing thresholds cause problems (measurement noise, numerics)

Exponentially decreasing thresholds with offset

Combine the advantages

Theorem

Consider system (1) with undirected and connected graph G. Suppose that

$$f_i(e_i(t)) = |e_i(t)| - (c_0 + c_1 e^{-at})$$

with $c_1, c_2 \ge 0$, at least one positive, and $0 < a < \lambda_2(G)$. Then, for all $x_0 \in \mathbb{R}^N$, the system does not exhibit Zeno behavior and for $t \to \infty$,

$$||\delta(t)|| \leq \frac{\lambda_N(G)}{\lambda_2(G)}\sqrt{N}c_0$$

Remarks:

- Size of final region is tunable
- Larger thresholds for small t increase inter-event times
- Problems due to noise or numerics are avoided

Next lecture

Guest lecture: Prof. Petter Ögren